

ERAMCO SYSTEMS

ES 84091A
MLDL ROM
annotated listing

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1785 LS Den Helder
The Netherlands

ERAMCO SYSTEMS
Kruiszwijn 2102
1788 LN Den Helder

```

*
*
*
* XROM number is 11
*
* Number of functions is 28
*
*      ORG          $000
*
*      FCB          $00B      XROM number is 11
*      FCB          $01E      number of functions is 28
*      FAT          $FDE      ERAMCO-MLDL
*      FAT          $4AE      RAMWR
*      FAT          $DBA      MMTORAM
*      FAT          $7A4      AFAT
*      FAT          $756      DFAT
*      FAT          $680      MOVE
*      FAT          $6FC      CLBL
*      FAT          $EE2      COPYR
*      FAT          $E5F      ROMSUM
*      FAT          $CAE      REG>ROM
*      FAT          $7DA      --
*      FAT          $316      COMPILE
*      FAT          $6D3      LOCA
*      FAT          $6B4      LROM
*      FAT          $6EC      COD
*      FAT          $6F6      DECOD
*      FAT          $F01      ROMCHKX
*      FAT          $C4B      ROM>REG
*      FAT          $803      MNEM
*      FAT          $E8B      DISASM
*      FAT          $44C      CAT
*      FAT          $715      CBT
*      FAT          $F88      SYNT
*      FAT          $F9D      GE
*      FAT          $7DA      --
*      FAT          $364      SAVEROM
*      FAT          $3BE      GETROM
*      FAT          $089      CMPDL
*      FAT          $047      IPAGE
*      FAT          $FE3      MKPR
*
*      FCB          $000
*      FCB          $000      end of FAT
*
*
*
*****
*
*      BEGIN OF ASSEMBLY LISTING
*
*****
*

```

```

*
*
* Initialize ram PAGE
*
      NAM      20 IPAGE    say we want a two digit numeric argument
*
      FCB      $000       say non programmable
      A<>B      X          save xrom number in "B(X)"
      GOSUB     VPAG       page number to "C[6]" and $400 to "C(X)"
      C=0       X          get rid of the loop counter
      R= 5
      A=C       A
      A=A-1     R<
      A<>C      M          end of page to "C(M)" start to "A(M)"
*
* here the page is completely cleared
*
IPAGE1  WRIT
      C=C-1     M
      ?A<C     M
      JC-       IPAGE1
*
* now we have to load the xrom number, the start address of the rom name
* and the number of functions
*
      C=B       X
      WRIT                      write the xrom number
      C=C+1     M
      LDI      $001
      WRIT                      say we have only a name
      C=C+1     M
      C=C+1     M
      LDI      #091
      WRIT                      load address of rom name
      A<>C      X
      RCR 3
      A<>C      X
      RCR 11
      LDI      $3E0
      WRIT                      rom name is terminated by a RTN
      M=C                      save address of name
*
* we convert the name into the lcd characters by getting the name out of
* alpha to the display. the name is left justified in the lcd
*
      ?NC XQ      ARGOUT  convert name
      ?NC XQ      ENLCD
      C=M
      C=C-1     M          create address of first character
      R= 13
      LD@R 10
      A<>C      A          load only 11 characters
*
* the rom name has to be written in inverted order. this is accomplished
* by getting the characters from the left of the display and write them to
* the ram page
*

```

```

IPAGE2  READ E          get leftmost character and left justify
        A<>C            M
        WRIT           write character of name
        C=C-1          M
        A<>C            M
        A=A-1          MS
        JNC-           IPAGE2  are we finished yet
                                no, do the rest
*
* the end of the rom name is flagged by bit 7 of the last character. we
* have to tell the mainframe where our rom name ends by setting this bit
*
        A<>C            M
        C=C+1          M
        FETCH           get last character
        C<>ST
        SETF 7
        C<>ST           flag the character
        WRIT
        ?NC XQ          ENCP00
        GOTO           EXIT  say "READY" and give address of first free word
*
*
*
* CoMPile and Delete Labels
*
        NAM            10 CMPDL  indicate an alpha prompt
*
        FCB            $000      say this program is non programmable
*
* indicate we want a pre compiled program by clearing flag 0,1 and 2
*
        C=0            A
        RAMSLCT
        READ E
        RCR 3
        R= 1
        C=0            R<      clear flags 0,1, and 2
        RCR 11
        WRIT E
*
* before compiling a program it is necessary to make sure that the program
* has a normal END. in case we have a program that is the last program in
* memory, we make the .END a normal one and load a new .END
*
COMPIL  GOSUB           FPAL
        RCR 8
        A<>C            R<      address of end of program to "A[3-0]"
        C=0            A
        RAMSLCT
        READ C          get address of .END
        ?A#C            X      is it the .END
        JC+             NO.END no, continue
        C=C-1          X      yes, create address of new .END
        C<>B            A      and save reg C in "B"
        ?NC XQ          AVAIL
        ?C#0            X      is there still an empty register under the .END
        ?NC GO          PACKE  no, pack and say "TRY AGAIN"

```



```

*
* we do have enough room to create the new .END this end has a link of one
* register to the old .END
*
      C=B          X
      RAMSLCT      address the new .END
      C=0          A
      R= 5
      LD@R C       say it is an end
      LDI          #120 link is 1 reg, make it .END, decompiled
      WRITDAT      new .END to memory
      C=B          X
      C=C+1        X
      RAMSLCT
      READDAT      get old .END
      R= 1
      LD@R 0       make it a normal END
      WRITDAT      END to memory
      C=0          A
      RAMSLCT
      C<>B         A
      WRIT C       store the updated reg C
*
*
*
NO.END  GOSUB      FPAL    find links
        CLR 10     say we are in ram
*
* FPAL returns from FLINK. the address of the end of the program is then
* returned in "C[11-8]". this is used by CPGMHD to find the start of the
* program we have to handle. next thing to do is to pack the program.
*
      RCR 8
      A=C          R<
      ?NC XQ       CPGMHD
      ?NC XQ       PUTPCF
      ?NC XQ       PACKN
*
* tell the user we are handling the 2 byte goto's
*
      ?NC XQ       CLLCDE
      ?NC XQ       MESSL
      RMB          $0A    "COMPL 2B G"
      ?NC XQ       LEFTJ
      ?NC XQ       ENCP00
      ?NC XQ       GETPC
*
* the main loop for the handling of the 2 byte goto's is entered with the
* address of a function in "A[3-0]" in MM format
*
DO2B    A<>C       R<
        A=C        R<    copy PC to "C[3-0]"
        N=C        and to "N"
        ?NC XQ      NXBYTA get the first byte of this instruction
*
* split the function byte into "A(MS)" and "A[0]". in this way we can save
* the number we have to jump to in "A(MS)" and test if we have a two byte
* GTO
*

```

```

C=0      XS
RCR 1
C=0      XS
A<>C     A
LDI      $00B
?A#C     X      do we have a two byte GTO
JNC+     F2BINS  yes, go handle it
C=N
A=C      R<
?NC XQ   NXLSST  get address of next instruction into "A[3-0]"
?FSET 6   have we reached the end of the program yet
JNC-     D02B    no, try the next instruction
JNC+     STPP3   yes, go handle the 3 byte GTO and XEQ
*
* here we have found the 2 byte instruction we have to compile. first thing
* we must do is to get the PC at the right position for calculating the
* distance to its label. therefore the PC has to be positioned to the last
* byte of the GTO
*
F2BINS   A<>B     A      save label number in "B(MS)"
C=N
A=C      R<
?NC XQ   INCAD2
?NC XQ   PUTPC      PC is positioned to the last byte of the GTO
CLRFB 8   say we are jumping forwards
C<>B     A
C=0      X
RCR 13    get label number to "C(X)"
C=C-1    X      decrement it by one for SEARCH
A=C      X      and put it in "A(X)" for SEARCH
*
* SEARCH returns the address of the wanted label in "C[3-0]" or if it is
* not found it will clear "C" also is the label number saved in "G". this
* is important, because INBYT expects the byte to be inserted in "G"
*
      GOSUB      FLABEL  find address of label, is label found
*
* FLABEL will return at the next instruction if the label is found, else
* it will skip this instruction
*
      JNC+       5      yes, go calculate the distance and direction
      GOTO      NOLBLF  no, error with "NO LBL NN" and position to GTO
*
* stepping stone
*
STPP4    JNC-     D02B
*
      M=C        save address of label in "M"
      ?NC XQ     GETPC  address last byte of GTO to "A[3-0]" + "C[3-0]"
      N=C        and save it in "N"
      C=M
*
* in the next part we test if we have to jump backwards or forward and we
* compute the distance in registers and bytes to the label
*

```

```

?A#C      X      are we in the same register
JC+        4      no, test if we are in a higher register
?A<C      R      is it a jump forward
JNC+        6      yes, we can just compute the distance
JNC+        3      go say we are jumping backwards
?A<C      X      is it a jump forward
JNC+        3      yes, we can just compute the distance
SETF 8      say we are jumping backwards
*
* when we jump backwards we have to make sure we get valid results when we
* compute the distance. this is done by exchanging the label and gto
* addresses because the label address is in fact in a higher numbered part
* of MM and distance is calculated by subtracting "C" from "A"
*
      A<>C      R<
      ?NC XQ      CALDSP go compute the distance
*
* when the distance is more then 15 register we have to convert the 2 byte
* GTO into a 3 byte GTO
*
      ?A#0      XS      is the distance more then 256 registers
      JC+        MK3BG  yes, go convert
      LSHFA      X      no, left justify number of registers
      ?A#0      XS      is it more then 15 registers
      JC+        MK3BG  yes, go convert
      LSHFA      X      no, left justify number of registers
      RSHFA      R<      put # reg and # bytes into "A(X)"
      C=N
      A<>C      R<      address of second byte to "A[3-0]"
      ?FSET 8      is the jump forward
      JNC+        2      yes, leave direction bit zero
      C=C+1      R      no, make direction bit one
      C=C+C      R<
      C=C+C      R<
      C=C+C      R<
      C=C+C      X
      RCR 2      left justify distance byte in "C[1-0]"
      ?NC XQ      PTBYTA and put it in the GTO
      JNC-      STPP4 do the next instruction
*
* stepping stone
*
STPP3      JNC+      D03BGX
*
*
* a three byte GTO is made by inserting one byte. the byte is inserted
* after the last byte of the two byte GTO. the byte to be inserted is
* expected to be found in "G". this byte has to have the number of the
* label we want to jump to. this is all accounted for, because the SEARCH
* routine transferred the byte to "G" and the address in "N" is of the last
* byte of the 2 byte GTO. also is the first byte of the 2 byte GTO replaced
* with $0D0, the identifier of a 3 byte GTO
*

```

```

MK3BG  C=N
      A=C      R<      get address of last byte to "A[3-0]"
      ?NC XQ    INBYT    insert the byte at next address
      R= 3
      ?NC XQ    DECADA
      ?NC XQ    DECADA    place the address at the first byte of the GTO
      LDI      $0D0
      ?NC XQ    PTBYTA    and make it a 3 byte GTO
      C=0      A
      RAMSLCT
      GOTO      COMPIL    find start of program, pack and start again
*
* in the next part the 3 byte GTO and XEQ are handled
*
* the address of the first instruction is saved and the message
* "COMPL 3B G/X" is placed in the display to keep the user informed of what
* is happening.
*
DO3BGX  ?NC XQ    PUTPC    save PC
      ?NC XQ    CLLCDE
      ?NC XQ    MESSL
      RMB      $0C      "COMPL 3B G/X"
      ?NC XQ    LEFTJ
      ?NC XQ    ENCP00    message is in display
      ?NC XQ    GETPC    get PC back to start
NXTINS  A<>C      R<
      A=C      R<
      N=C      copy current address to "A[3-0]" and "N"
*
* in this part we get the first byte of an instruction out and are testing
* if it is a 3 byte GTO or a XEQ. if so, we are compiling it
*
      ?NC XQ    NXBYTA    get first byte of instruction
      C=0      XS
      RCR 13
      C=C+1    XS
      JC+      ROWF      yes, go do the next instruction
      C=C+1    XS
      JC+      DOCOMP     yes, go compile it
      C=C+1    XS
      JC+      DOCOMP     yes, go compile it
*
* we have no GTO or XEQ, so we just find the next instruction and loop
* if we haven't reached the end of the program yet
*
ROWF    C=N
      A=C      R<      get current address to "A[3-0]"
      ?NC XQ    NXLSST    find next instruction
      ?FSET 6
      STPP2    JNC-      NXTINS    no, do next instruction
      GOTO      ?READY    yes, decide whether we can exit or do more
*
* here we are first trying to find the desired label. the label number is
* given in the third byte of the instruction. also we find a direction bit
* here. this bit must be cleared because if it is set, we would search for
* a label number that is 128 to high and will always find an error then
*

```

```

DOCOMP  ?NC XQ      PUTPC  save address of first byte in PC
        SETF 8      say jump is forward
        ?NC XQ      GT3DBT byte with label number in "ST" address in "M"
        CLRf 7      clear direction bit
        C<>ST
        A=C          X      byte to "A(X)" for SEARC1
        SETF 6      say we are at first byte of a 3 byte instr.
*
* FLABEL returns the address of the label if it exists. if it exists, it
* will execute the next step, else the next step is skipped
*
        GOSUB        FLABEL  find label address, have we found it
        JNC+         LBLF
        GOTO         NOLBL   no, go error "NO LBL NN"
*
* stepping stone
*
STPP1   JNC-         STPP2
*
*
* in this part we find out in which direction we have to jump.
*
LBLF    A=C          R<      address of label to "A[3-0]"
        C<>M          address of GTO to "C[3-0]"
        ?A#C          X      are we in the same register
        JC+           4      no, test if we are in a higher register
        ?A<C          R      are we jumping backwards
        JNC+          DIST   no, go calculate the distance
        JNC+          3      yes, say it and fix addresses for calculation
        ?A<C          X      are we jumping backwards
        JNC+          DIST   no, go calculate the distance
        CLRf 8        say we jump backwards
        A<>C          R<      and fix the addresses for the calculation
*
* here we are computing the distance in registers and bytes
*
DIST    C=A-C          X      calculate number of registers
        C=A-C          R      calculate number of bytes
        JNC+           4
        C=C-1          X      if we pass register boundary decrement reg 1
        C=C-1          R
        C=C-1          R      and set byte count to right value
*
* now we have found the distance in registers and bytes. in bit 3-1 of
* digit "C[3]" we have the byte distance. bit 0 is used to have one bit
* of the register count. this count is placed in bit 0 of "C[3]" and
* "C[2-1]".
*

```

| | | | |
|---|--------|---|--|
| C=C+C | X | | |
| C=C+C | X | | |
| C=C+C | X | | |
| C=C+C | X | | left justify reg count and use bit 0 if needed |
| JNC+ | 2 | | |
| C=C+1 | R | | we have to use bit 0 |
| RSHFC | R< | | place distance digits in "C(X)" |
| C<>B | X | | and save them in "B(X)" |
| ?NC XQ | GETPC | | get address of first byte |
| ?NC XQ | GTBYTA | | and get the byte |
| RCR 12 | | | place row digit in "C[3]" |
| C=B | X | | add the distance digits |
| RCR 2 | | | leave first byte in "C[1-0]" |
| C<>B | M | | |
| C<>B | MS | | save second byte in "B[13-12]" |
| ?NC XQ | PTBYTA | | put first byte in memory |
| ?NC XQ | INCADA | | |
| C=B | A | | |
| RCR 12 | | | get second byte |
| ?NC XQ | PTBYTA | | put second byte in memory |
| ?NC XQ | NXBYTA | | get last byte |
| C<>ST | | | |
| CLRF 7 | | | |
| ?FSET 8 | | | are we jumping backwards |
| JNC+ | 2 | | |
| SETF 7 | | | yes, indicate a jump backwards |
| C<>ST | | | |
| ?NC XQ | PTBYTA | | put byte in memory |
| JNC- | STPP1 | | go do the next instruction |
| * | | | |
| ?READY | C=0 | A | |
| RAMSLCT | | | |
| READ E | | | get flags |
| RCR 3 | | | |
| C<>ST | | | |
| ?FSET 0 | | | do we have to do a precompiled program |
| JNC+ | 4 | | |
| GOTO | READY | | no, do exit or return to calling program |
| SETF 0 | | | say we do compile for a second run |
| SETF 1 | | | say we use the user regs as label addresses |
| C<>ST | | | |
| RCR 11 | | | |
| WRIT E | | | save flags |
| * | | | |
| ?NC XQ | GETPC | | |
| ?NC XQ | FLINK | | |
| RCR 8 | | | |
| A=C | R< | | |
| ?NC XQ | CPGMHD | | |
| ?NC XQ | PUTPC | | position PC to start of program |
| * | | | |
| * here the loop for finding GTO or XEQ indirect functions start. we fetch | | | |
| * the first byte fo every function and testif it is IND. if so, we exit | | | |
| * | | | |

```

FNDIND  ?NC XQ      NXBYTA
        A<>B      R<      save current address in "B[3-0]"
        A<>C      X
        A=0      XS      instr. byte to "A(X)"
        LDI      $OAE
        ?A#C      X      is it XEQ/GTO indirect
        JC+      4      no
        GOTO      X/GIND  yes, position PC to this line and exit
        A<>B      R<
        ?NC XQ      DECADA  address to a standard position
        ?NC XQ      NXLSST  find next instruction
        ?FSET 6      are we finished yet
        JNC-      FNDIND  no, test next instruction

```

*
* the addresses of the labels are stored in the user registers. therefore
* it is needed to clear these registers and to find the address of the
* first register. this address is stored in reg 10 that is used to keep
* track of the next free register.
*

```

        C=0      A
        RAMSLCT
        READ C
        RCR 3      get address of reg 0
        A<>C      X
        READ 10
        A<>C      X
        WRIT 10      put it in reg 10
        SETF 8
        ?NC XQ      CLR      clear the user registers

```

*
* in this loop the labels are searched, and if one is found the label is
* cleared and its address is saved in the user registers and its number.
*

```

STDLBL  ?NC XQ      GETPC  get start address of program
        M=C      save it in "M"
DLLBL1  ?NC XQ      NXBYTA  get first byte of an instr.
        A<>B      R<      save current address
        R= 1
        C=C-1      R      is it row 0
        JC+      LBL014  yes, go handle row 0
        A=C      X
        LDI      $0BF      first byte of 2 byte LBL - $10
        ?A#C      R<      is it a 2 byte label
        JNC+      LBL>14  yes, go handle it
DLLBL2  R= 3
        A<>B      R<      get current address back to "A[3-0]"
        ?NC XQ      DECADA  to standard position
        ?NC XQ      NXLSST  find next instr.
        ?NC XQ      PUTPC  save address in PC in case we have to pack
        ?FSET 6      have we reached the end yet
        JNC-      DLLBL1  no, do next instr.
        GOTO      NO.END  yes, go compile the program

```

*
* in this part we separate the numeric labels from the local labels. in
* case we have a local label we must not delete it. it is necessary to
* store the address of the label for the label find routine.
*

```

LBL>14  R= 3
        A<>B    R<
        ?NC XQ      NXBYTA  get label number
        A<>B    R<          save address
        A=C        X
        A=0        XS          label number to "A(X)"
        LDI        $064
        ?A<C      X          is it a local label
        JC+        3          no
        A<>B    R<
        JNC-      DLLBL1
        R= 0
        A<>C      X
        C=C+1    X          increment LBL for it is decremented when stored
        G=C          save label number in "G"
        R= 3
        A<>B    R<          get address of second byte
        C=0      A
        ?NC XQ      PTBYTA  delete the second byte
        ?NC XQ      DECADA  point to first byte
        JNC+      SECBYT

```

*
 * here is checked if we have to do with a null byte, and if so we return to
 * the main loop
 *

```

LBL014  LD@R 0          clear the first digit and set pointer to 0
        ?C#0      R          is it a null byte
        JNC-      DLLBL2  yes, go do the next instr.
        G=C          save label number in "G"
        R= 3
        A<>B    R<          current address to "A[3-0]"

```

*
 SECBYT C=0 A
 ?NC XQ PTBYTA delete the byte
 *

* here we are combining the start address of the program, the address of
 * the label and its number
 *

```

        R= 3
        ?NC XQ      GETPC  get the address of the label to "A[3-0]"
        C=M          get start address of program
        RCR 10
        A<>C    R<          add label address
        RCR 11
        R= 0
        C=G          add label number
        C=C-1      X
        C=0        XS
        C=0        MS
        C=C+1      MS          type it as alpha data
        A<>C      A
        C=0      A
        RAMSLCT
        READ 10          get storage address
        C=C+1      X
        WRIT 10          update storage address
        C=C-1      X          create the right storage address
        RAMSLCT
        A<>C      A
        WRITDAT          put the label address in the register

```


*
 * after we have cleared the label we have to pack the program again and
 * start from the beginning. this is necessary, because we have created a
 * null byte in the memory. we have to change the status of the program from
 * packed into unpacked for PACKN skipps packed programs.

```
*
      ?NC XQ      GETPC
      ?NC XQ      FLINK      get address of END to "C[11-8]"
      ?NC XQ      FIXEND     make the program unpacked
      ?NC XQ      PACKN
      GOSUB        FPAL
      RCR 8
      A<>C      R<
      ?NC XQ      CPGMHD
      ?NC XQ      PUTPC      start address of program to PC
      GOTO        STDLBL

*
*
*
NOLBL  ?NC XQ      GETPC      set PC to the 2 byte GTO that has no LBL
      ?NC XQ      DECADA
      JNC+        3
NOLBLF ?NC XQ      GETPC      set the PC to the GTO/XEQ that has no LBL
      ?NC XQ      DECADA
      ?NC XQ      PUTPC
      ?NC XQ      CLLCDE
      ?NC XQ      MESSL
      RMB         $07      "NO LBL_" put this in the display
      R= 0
      C=6
      A=C         X          get label number to "A(X)"
      A=0         MS
      A=0         XS
      A=A+1       MS
      A=A+1       MS
      ?NC XQ      GENNUM      say we want two digits in the display
      ?NC XQ      LEFTJ      append label number to the display
      ?NC XQ      ENCP00
      ?NC XQ      MSGDLY
      ?NC GO      ERR110     put total message in the display
```

*
 * here we indicate the user that we are ready and place the PC at the first
 * step of the compiled program

```
*
READY  ?NC XQ      CLLCDE
      ?NC XQ      MESSL
      RMB         $05      "READY"
      ?NC XQ      LEFTJ
      ?NC XQ      ENCP00
      ?NC XQ      STMSGF     message is in display
```

*

```

GOSUB      FPAL
RCR 8
A=C      R<
?NC XQ    CPGMHD  get start of program to "A[3-0]"
?NC XQ    PUTPCF  put user at start of program
READ E
RCR 3
C<>ST
?FSET 2
?NC GO    TONE7X  give audio warning that we finished and exit
READ P
RCR 7
GOTO ADR

*
*
*
* Find Program And Links
*
* Entry: register Q holds the inverted right justified name of the program
*        you want to have the begin and end address from
* Exit : "C" has links to the desired label as returned by FLINK
* Uses : all
*
FPAL      R= 3
          READ Q
          M=C      get a copy of the program name to "M"
          ?C#0     R<  do we have to do the current program
          JNC+     CURPG
          ?NC XQ   ASRCH  no, find program
          ?C#0     R<
NONEX     ?NC GO   ERRNE  say program is nonexistent
          ?FSET 9
          JC-      NONEX  it is a microcode function
          ?FSET 2  is the program hold in main memory
          JNC+     PFOUND yes
ROM       ?NC XQ   ERROR
          FCB      $06A  specify message 'ROM'
CURPG     ?FSET 10 are we in ROM at the moment
          JC-      ROM   yes
          ?NC XQ   GETPC
PFOUND    ?NC GO   FLINK  find program head and it's END *
*
*
*
X/GIND    A<>B     R<  get address of first byte of GTO/XEQ ind
          ?NC XQ   DECADA
          ?NC XQ   PUTPC  set PC to the error line
          ?NC XQ   CLLCDE
          ?NC XQ   MESSL
          RMB      $0B    "GTO/XEQ IND" message to display
          ?NC XQ   LEFTJ
          ?NC XQ   ENCP00
          ?NC XQ   STMSGF
          LDI      $004
          ?NC GO   TONEB  give low warning beep if audio is enabled
*
*
*

```

```

*
* this is the label find routine. we decide here if we have to do with
* the normal compiler or if we have to use the user registers with the
* addresses of the labels and their label number. it depends on flag 1 in
* reg E which compiler we are doing
*
FLABEL  A=0      XS
        LDI      $064
        ?A<C      X
        JNC+      LOCLBL
        C=0      A
        RAMSLCT
        READ E
        RCR 3
        C<>ST      get compiler flags
        ?FSET 1      do we have to use the user regs
        JC+      DUSREG  yes, go handle the user regs
*
* it depends on the setting of flag 6 of the original flag status if we
* have to do 2 or 3 byte GTO/XEQ
*
        C<>ST      get old flagstatus back
        ?FSET 6      is it 3 byte instr.
        JC+      4      yes, do three byte instr.
        ?NC XQ      SEARCH  no, do 2 byte
        JNC+      3
LOCLBL  ?NC XQ      SEARC1  do 3 byte
*
* if we have found the label we will do a normal return to the instruction
* after the calling subroutine, else we skipp this instruction and do the
* next instruction.
*
        ?C#0      A      does the desired label exist
        ?C RTN      yes, do a normal return
        POP      no, skipp instruction after the call
        C=C+1      M
        GOTOADR
*
* for a correct compilation of the program we must start searching for the
* desired label from our current program position. this makes it neccessary
* to search for the register with the address after our current position.
* when this register is found, we can search for the desired label. when we
* have reached the end of the labels and haven't found the desired label
* yet, we must start searching from the start of the program. this means we
* have to start with the first user register.
*
* this is also the procedure that is followed by the normal search routine
* in the mainframe
*
DUSREG  C<>ST      restore flags for compile
        A<>C      X      get the desired label number
        R= 0
        G=C      and store it in "G"
*
* in this loop we search for a label with an address that lays behind the
* current position in the program. we exit when we have found such a label
* or when we have reached the end of the labels.
*

```

```

R= 3
?NC XQ      GETPC    current address to "A[3-0]"
READ C
RCR 3              address of first reg to "C(X)"
*
FLADD  RAMSLCT      select a label register
      C<>B          X    save address counter in "B(X)"
      READDAT
      RCR 3          get address right justified
      ?C#0          A    have we had all labels
      JNC+          FLBL  yes, go find the desired label
*
* in the next instructions we test if the label is after the current
* position in the program.
*
      ?A#C          X    same register
      JC+          FLBL  no, test if higher reg
      ?A<C          R<    is address of label higher
      JNC+          6    yes, go search for desired label
      C<>B          X    we have not found such a label yet and are
      C=C+1          X    going to do the next label
      JNC-          FLADD go try next register
      ?A<C          X    is address of label higher
      JC-           4    no, do next label
*
* after we have found where we must start searching for the labels we will
* do it here. if the desired label is not found yet when we reach the end
* of the labels we will start at the beginning of the labels.
*
FLBL   R= 0
      C=G
      A=C          X
      A=0          XS    desired label number to "A(X)"
      C<>B          X
*
DNLBL  RAMSLCT
      C<>B          X    save label counter
      READDAT      get a label
      ?C#0          A    have we had all labels
      JNC+          WRAP  yes, start at the beginning of the labels
      ?A#C          X    is this the desired label
      JNC+          LBLFND yes, exit
      C<>B          X
      C=C+1          X    point to next label
      JNC-          DNLBL do next label
*
WRAP   C=0          A
      RAMSLCT
      READ C
      RCR 3          get counter for first address and go search
      JNC-          DNLBL
*
* we have to exit the same as the SEARCH and SEARC1 routine. this means
* status registers must be selected and address of label is in "C[3-0]"
* and label number is in "G"
*

```

```

LBLFND  RCR 3          address to right position
        R= 3
        A=C          R<    save in "A" while we select status regs
        C=0          A
        RAMSLCT      select status regs
        A<>C        R<    address of label back to "C[3-0]"
        RTN          return to calling point
*
*
*
* COMPILE
*
        NAM          10 COMPILE indicate an alpha prompt
*
        FCB          $000    say non programmable
        C=0          A
        RAMSLCT
        READ E
        RCR 3
        ST=0
        SETF 0          say we want the normal compile routine
        C=ST
        RCR 11
        WRIT E          save flag status for compile
        GOTO          COMPIL do it
*
*
*
* this routine checks if the HPIL rom is present
*
PILTST  C=0          A
        R= 6
        LD@R 7
        FETCH          get xrom number of page 7
        A=C          X
        LDI          $01C
        ?A#C          X    is it the HPIL chip
        ?NC RTN        yes, then return to calling routine
        GOSUB          INERR no, initialize for error message
        RMB          $07    "NO HPIL"
*
ERREX   ?NC XQ          LEFTJ  exit error message and return to mainframe
        SETF 8
        ?NC XQ          MSG105
        ?NC GO          ERR110
*
*
*
* this routine checks if the page address is valid and it sets the loop
* counter to $400. in the loop we store 4 rom words, so we do a complete
* rom block of 4K.
*

```

```

VPAG    READ X
        ?NC XQ      BCDBIN
        A<>C        X      get wanted page to "A(X)"
        LDI          $010
        ?A<C        X      is it a valid page address
        JC+          VPG
        GOSUB        INERR   no, put message to display and exit
        RMB          $09     "PAGE > 15"
        JNC-         ERREX

*
*
*
* here is the initialization done for an error message
*
INERR    ?NC XQ      ERRSUB
        ?NC XQ      CLLCDE
        ?NC GO      MESSL

*
*
*
* SAVEROM
*
        NAM          SAVEROM

*
        GOSUB        PILTST
        GOSUB        VPAG
        C=0          MS
        ?NC XQ        780B   check if there is already afile with this name
        C=M
        ?C#0          A      did the file exist
        ?C GO         7692   yes, say DUP FL NAME and exit to mainframe

*
* the romfiles do have an file type of 7. this is done so we can't destroy
* it with any of the normal HPIL functions. the length is 640 registers and
* these register are stored in 20 records. the file is automatically
* secured when it is created.
*
        C=0          A
        LDI          $078   file type 7 and secured
        RCR 2
        LDI          $280   640 registers
        RCR 4
        LDI          $014   20 records
        RCR 4
        ?NC XQ        76BC   create the file
        ?NC XQ        7F72   set tape on start of file and BP to zero

*
        LDI          $0A2
        ?NC XQ        70BA   give DDL 2 - write mode

*
        ?NC XQ        70DA   give DDL 0 - write buffer 0

*
        ?NC XQ        77E7   do error check

*

```

```

        SELP 4
        FCB          $005      set REG 1 for DAB
*
        GOSUB        VPAG      get rom page and loop counter to "C"
*
* the main loop of SAVEROM is started here. we store 5 bytes after each
* other and these bytes are containing the four lower bytes and one upper
* byte. they are stored sequentially. the first byte is the upper byte
*
MLSR    C=C-1      X          have we had the whole 4K block
        JC+        CLUP      yes, clean up and exit
        R= 4              set loop counter to do 4 romwords
        A=0          X
        C<>B        X          save main loop counter in "B(X)"
*
* in this loop we make five bytes containing four rom words
*
D4RWRD  FETCH              get a rom word
        C=C+1      M
        A<>C        XS
        A<>C        X          lower part to "A(X)", high PART TO "C(XS)"
        LSHFA      X
        LSHFA      A
        LSHFA      A          right justify low bytes in "A(M)"
        C=C+C      X
        C=C+C      X
        RSHFC      X
        A<>C        X          left justify high bits in "A[1-0]"
        R=R-1
        ?R= 0              have we had 4 romwords
        JNC-          D4RWRD no, do next word
*
        C<>B        X
        M=C              save mainloopcounter and current address
        LSHFA      X
        LSHFA      A          right justify the 5 bytes in "A[11-2]"
        A<>C        A
        R= 5              set loop counter for 5 bytes to be saved
        WRIT Q            string bytes to reg Q
*
* it is neccessary to use register Q because we can't use "N" for this
* holds some essential file parameters
*
S5BYT   READ Q
        RCR 2
        WRIT Q              get one byte to "A[1-0]"
        ?NC XQ          7126 write it on tape
        ?NC XQ          77E7 test for transmit error and exit if error
        R=R-1
        ?R= 0              have we written 5 bytes
        JNC-          S5BYT no, do next byte
        C=M              get loop counter and current address
        JNC-          MLSR do next words
*
* here the last record is closed and the cassette drive retrieves its
* normal status and we exit to mainframe
*

```

```

CLUP      LDI          $0AB
          NOP          (0C0)    this is done to solve a bug, these 4
                                words could also be deleted

          NOP
          ?NC GO        70AF

*
*
*
* GETROM
*
          NAM          GETROM

*
          GOSUB        PILTST   error if there is no HPIL chip
          GOSUB        VPAG     test for valid page
          LDI          $007
          ?NC XQ        780A    error if this is not our file type

*
          ?NC XQ        7F77    seek the file
          ?NC XQ        70E6    send buffer 0 and SDA
          GOSUB        VPAG     find rampage to write at and set loop counter

*
* the mainloop of getrom starts here. we get 5 bytes out of the tape after
* each other and combine these into 4 romwords
*
MLGR      C=C-1      X
          ?C GO        70AC
          M=C
          R= 5
                                have we had the whole 4K
                                yes, UNT and return to mainframe
                                save loop counter and current address
                                set counter for getting 5 bytes

*
* here we are getting 5 bytes out and save them in reg Q
*
G5BYT     ?NC XQ      7110    get a byte
          ?FSET 9
          ?C GO        7634    did we have any type of error
                                yes, decide which error and say it
                                send byte out again
          WRITHPIL
          A<>C      X
          READ Q
          A<>C      X
          RCR 2
          WRIT Q
                                append byte to string in reg Q
          R=R-1
          ?R= 0
                                have we done 5 bytes
          JNC-        G5BYT    no, do next byte

*
* after we have found 5 bytes, we have to store them in the desired page
*
          A=C          A
                                save the 4 lower bytes left justified in "A"
          RCR 2
                                get upper byte to "C[3-2]"
          C<>B      A
          C=M
                                get current address and main loop counter
          R= 4
                                set counter for 4 romwords
          C<>B      A
                                initialize for loop to store 4 words

*
* the loop has to be entered with the 4 lower bytes left justified in "A"
* and the upper byte right justified in "C[3-2]"
*

```


| | | | |
|-------|--------|-------|--|
| D5BYT | A<>C | A | |
| | RCR 12 | | |
| | A<>C | A | shift next lower byte to "A[1-0]" |
| | A<>C | XS | append the upper bits to the word |
| | A<>C | X | word to "C(X)" |
| | C<>B | M | get destination address |
| | WRIT | | word to page |
| | C=C+1 | M | increment destination address |
| | C<>B | M | dest address to "B(M)" |
| | C=C+C | A | |
| | C=C+C | A | |
| | RCR 1 | | shift upper byte two bits to the right |
| | R=R-1 | | |
| | ?R= 0 | | have we done 4 words |
| | JNC- | D5BYT | no, do next word |
| | C<>B | A | get loop counter and current address |
| | JNC- | MLGR | do next 4 words |

```

*
*
*
* CAT and RAMWR routines
*
*
*
* APPend Hex Digit
*
* after a key has been pushed during partial key sequence, the main program
* may call this subroutine for processing the key. only hex digits are
* accepted ( i.e. digits 0-9 characters A-F ).
*
* Entry: data as produced by the HP-41 operating system after a key has
*        been pushed. the display should be enabled on entry
* Exit  : if the key is digit 0-9 or chracter A-F the corresponding
*        character is appended to the display. the address after the call
*        is skipped. for all other keys 70 ms visual feedback. the next
*        instruction is not skipped
* Uses  : "A(MS)", "A(X)", "C", "N" and the active pointer
*
APPHD      ?FSET 4          APP1      upper 10 keys
          JC+              APP1      yes
          ?FSET 3          APP1      is it a digit key
APP2      ?NC GO          BLINK      no, give 70 ms feedback
          A<>C            MS         recall numerical value of the key
          RCR 13
          R= 1
          LD@R 3          APP3      make it a display character
APP3      WRIT E          APP3      append it to the display
          POP
          C=C+1           M
          GOTOADR         APP3      skipp byte after the call

```

| | | | | |
|------|--------|----|------|--|
| APP1 | ?A#0 | MS | | is it the TAN key |
| | JNC- | | APP2 | yes, it is not allowed |
| | R= 13 | | | |
| | LD@R 7 | | | |
| | ?A<C | MS | | is it G-J |
| | JNC- | | APP2 | yes, it is not allowed |
| | R= 2 | | | |
| | LD@R 5 | | | |
| | A=C | XS | | |
| | C=N | | | |
| | C=C+A | XS | | add 5 to keycode |
| | RCR 12 | | | |
| | R= 6 | | | |
| | LD@R 1 | | | |
| | LD@R 5 | | | load address of character in defaultcode table |
| | FETCH | | | get characte |
| | R= 1 | | | |
| | C=0 | R | | convert it to display character |
| | JNC- | | APP3 | add character to display and exit |

*
 *
 *
 * NULL
 *
 * waits one second for key release. if key not released after this period
 * "NULL" is displayed for at least ,3 seconds. when the key is released
 * within 1 second, the byte after the subroutine call is skipped
 *
 * Entry: appropriate text in the display (display is left justified by
 * NULL). assumes display is enabled on entry.
 * Exit : display is still enabled if the key is released in time, otherwise
 * chip 0 is enabled
 * Uses : "A", "C", active pointer, sethex, flag 8
 *

| | | | | |
|-------|---------|---|-------|-------------------------------------|
| NULL | ?NC XQ | | LEFTJ | |
| | LDI | | \$2B3 | |
| | C=C+C | X | | load timing constant for one second |
| NULL1 | CLRKEY | | | |
| | ?KEY | | | wait one second for key release |
| | JNC+ | | NULL3 | key is released in time |
| | C=C-1 | X | | |
| | JNC- | | NULL1 | |
| | CLRF 8 | | | say no blinking |
| | ?NC XQ | | MSGA | |
| | FCB | | \$03C | "null" |
| | LDI | | \$3E8 | load timing constant for .3 seconds |
| NULL2 | C=C-1 | X | | |
| | JNC- | | NULL2 | show message for .3 seconds |
| | ?NC GO | | RSTKB | wait for key release |
| NULL3 | POP | | | |
| | C=C+1 | M | | |
| | GOTOADR | | | skip byte after the call |

```

*
*
* Read HEX Digit 2
*
* reads specified number of hex digits from the display and returns these
* digits in "A(M)" right justified
*
* Entry: number of digits minus one is specified in "A(MS)" display must
*         be enabled on entry
* Exit : the wanted digits are returned in "A(M)" display is still
*         enabled
* Uses : "A", "C", "B(MS)", active pointer
*
* RHEXD1 reads 4 digits from the display
*
RHEXD1  R= 13
        LD@R 3
        A=C      MS      say we want four digits
RHEXD2  B=A      MS      save number of wanted digits in "B(MS)"
        A=0      M      initialize "A(M)" for holding digit string
        LDI      $009
        A=C      X      conversion constant for dig A-F to "A(X)"
RHEX 1  READ D    move digits to be read to the left of display
        A=A-1    MS      have we had all digits
        JNC-     RHEX 1  no, do next
        A<>B     MS      recall counter
RHEX 2  R= 1
        READ E    read one digit
        ?C#0     R      is it 9-0
        JC+      RHEX 3  yes
        C=C+A    R<     no, convert to A-F
RHEX 3  C=0      R
        RCR 11
        LSHFA    M      make place for new digit
        R= 3
        A=C      R      append digit
        A=A-1    MS      have we had all
        JNC-     RHEX 2  no, do next digit
        RTN
*
*
*
* CATalog
*
* ATTENTION !! this routine does not work in the expected way with the HP-
* 41 CX. this is caused by the fact, that the CAT routine from the 41CX is
* moved to page three with partial handling in the normal mainframe roms.
* however, this is not so important, because the 41CX already has a sort of
* port selectable catalog.
*
        NAM      CAT
*
        FCB      $000    say non programmable
        FCB      $000    say direct executing
*
        A=0      X
        ?NC XQ    CLLCDE
        ?NC XQ    PROMFC  display "CAT "
        JNC+     CATAL2

```

```

*
* when the user pressed backarrow we return here and cancel the function
* after we have done some housekeeping
*
CANCEL  ?NC XQ      CLLCDE
        ?NC XQ      ANNOUT
        ?NC XQ      RSTSEQ
        ?NC GO      NFRKB

*
* append a prompt sign and wait in light sleep for a key to be pressed
* if the key is a backarrow we return at byte after call, otherwise this
* byte is skipped
*
INVASK  ?NC XQ      NEXT1
        JNC-        CANCEL
        GOSUB        APPHD    append the hex digit to the display
        JNC-        INVASK    invalid digit, ask again
        A=0          MS       say we want to read one digit from display
        GOSUB        RHEXD2    get digit to "A(M)"
        GOSUB        NULL      wait for key release
        JNC-        CANCEL    if key hold down to long, cancel
        ?NC XQ      ENCP00
        ?NC XQ      RSTSEQ    clear user flags
        R= 3
        C=0          A
        A<>C        R          get CAT # to "C[3]"
        C=C+1        X
        RCR 3
        ST=C          save CAT # in "ST" ( needed by 0B84 )
        RCR 8
        A=C          M          address of second word at page to "A(M)"
        R= 6
        LD@R 5
        R= 6
        ?A<C        R          is it a regular CAT ( 1, 2 or 3 )
        ?C GO        OB84      yes, go do the CAT
        A=0          X
CAT 2    ?A#C        R          have we had all blocks before this page
        JNC+        CAT 1      yes
        FETCH
        A=A+C        X          add number of functions on page to total
        C=C+1        R          point to next page
        JNC-        CAT 2      do next page
CAT 1    A<>C        X
        RCR 4          number of functions to "C[12-10]" page address
        A=C          A          to "A(XS)" + "C(XS)" + functions to "A[12-10]"
        ST=0
        SETF 1
        ?NC GO        OB85      do rest of initialisation for 0B85
                                fall in the regular CAT 2

*
*
* Display ADdRess
*
* this routine takes a hexadecimal digit from reg Q and writes it to the
* display ( right justified ). a space is appended.
*
* Entry: none
* Exit : display is enabled
* Uses : "A", "C", and active pointer
*

```

```

DPADR  ?NC XQ      ENCP00
      READ Q
      R= 13
      LD@R 3
      A=C          A
      ?NC XQ      CLLCDE
DPBYT  LDI          $009  load comparing constant
      A=C          X
      A<>C        M
      RCR 7
      digits to display to "C[13-10]"
*
* in the loop is tested if the character is A-F or a normal digit. in case
* of the A-F we only have to subtract 9 to get the LCD character. else we
* must add $30 to the digit to get the LCD character.
*
DPADR2  RCR 13
      R= 1
      C=0          R
      ?A<C        R<
      JNC+        5
      A<>C        R<
      A=A-C        R<
      A<>C        R<
      JNC+        2
      LD@R 3
      WRIT C
      A=A-1        MS
      JNC-        DPADR2
      LDI          $020
      WRIT C
      RTN
      append character to the display
      have we had all characters
      no, do next one
      append a space to the string
*
*
*
* RAMWrite
*
      NAM          RAMWR
*
      FCB          $000  say nonprogrammable
*
ASKAD  ?NC XQ      CLLCDE
      ?NC XQ      MESSL
      RMB          $04  "ADR "
      JNC+        9
*
EXRMWR ?NC XQ      CLLCDE
      ?NC XQ      ANNOUT
      ?NC XQ      RSTSEQ
      ?NC GO      NFRKB
      update the display to the current status
      clean up the partial key sequence
      exit to mainframe
*
P4D    LDI          01F
      WRIT E
      ?NC XQ      NEXT3
      JNC-        EXRMWR
      append one prompt sign
      append 3 prompt signs and wait for key pressed
      it was backarrow, so exit
*
      GOSUB        APPHD
      JNC-        P4D
      no hex digit, so try again
*
P3D    ?NC XQ      NEXT3
      JNC+        6
      append 3 prompt signs and wait for key pressed
      it was backarrow
*

```

```

        GOSUB      APPHD
        JNC-       P3D      no hex digit, so try again
*
        JNC+       3        it was a hex digit, so do next digit
*
* the backarrow key was hit, so we remove the rightmost character and try
* again
*
        READ D
        JNC-       P4D      remove digit
                           ask for four digits again
*
P2D      ?NC XQ     NEXT2    append 2 prompts and wait for key
        JNC+       6        process backarrow key
*
        GOSUB      APPHD
        JNC-       P2D      no hex digit, so try again
*
        JNC+       3        skip removing part
*
        READ D
        JNC-       P3D      remove one digit and prompt for 3 again
*
P1D      ?NC XQ     NEXT1    append 1 prompt and wait for key
        JNC+       6        process backarrow key
*
        GOSUB      APPHD
        JNC-       P1D      no hex digit, so try again
*
        JNC+       3        skip removing part
*
        READ D
        JNC-       P2D      remove one digit and prompt for 2 again
*
        GOSUB      RHEXD1   get address to "A(M)"
        GOSUB      NULL     wait for key release
ASKAD1   JNC-       ASKAD    if nulled, ask for address again
        ?NC XQ     ENCP00
ANDW     A<>C       M
        A=C        M
        WRIT Q      save address in reg Q
        R= 5
        C=0        R<      create address of first word
        ?A#C       M        are we at first word
        JNC+       NFWRD    yes, continue
        FETCH
        ?C#0       X        is xrom zero
        JC+        NFWRD    no, continue
*
* when we have a first word that is zero, we must assume that there is no
* rampage. then we say "NO ROM" and go asking for the address again.
*

```

```

?NC XQ      CLLCDE
?NC XQ      MESSL
RMB          $06      "NO ROM"
ADRAG ?NC XQ      LEFTJ
?NC XQ      BLINK
LDI          $3E8
C=C-1      X          leave message .3 sec in display
JNC-        1
JNC-        ASKAD1

*
NFWRD GOSUB      DPADR      put the 4 hex address digits in the display
?NC XQ      ENCP00
?NC XQ      OFSHFT
READ Q
FETCH          get data at this word
RCR 10
R= 13
LD@R 2          say we want to append 3 digits
A=C      A
?NC XQ      ENLCD
GOSUB      DPBYT      PRINTDISPLAY

*
* the display now contains the address and the data at this address. now
* we are waiting for the user to press a key.
*
WFKEY ?NC XQ      NEXT      wait for key
JNC-   ASKAD1      if backarrow, ask for the address again
JNC+   2

*
* stepping stone
*
STP1   JNC-        NFWRD
*
C=N
RCR 1
C=0      XS
A=C      X          keycode to "A(X)"

*
* in this part we check for some special keys. these are SHIFT, BST, STO,
* SST and TAN. the keycodes that are returned are the synthetic keycodes
* minus one
*
LDI          $002
?A#C      X          is it shift
JNC+       TSHFT      yes, go set shift and wait for next key
LDI          $00A
?A#C      X          is it shift/shift
JC+        TSTK        no, test for other keys
TSHFT ?NC XQ      TOGSHF
?NC XQ      ANNOUT
?NC XQ      ENLCD
JNC-       WFKEY

*
TSTK LDI          $022
?A#C      X          is it STO
JC+        4
GOTO      RMWSTO      yes, go display contents of rom
*

```

```

        LDI          $041
        ?A#C        X          is it the TAN key
        JC+          4
DOBSTP  GOTO        RWRBST     yes, do a backstep
*
        C=C+1      X
        ?A#C        X          is it the SST key
        JC+          4
        GOTO        RMWSST     yes, do a single step
*
        LDI          $04A
        ?A#C        X          is it shift SST
        JNC-        DOBSTP     yes, go do a backstep
*
        ?FSET 3
        JC+          4          yes
BDD+A   ?NC XQ      BLINK      no
STP2    JNC-        STP1       display address and data again
*
* if we have to do with a digit key, we remove the data digits
*
        READ D
        READ D
        READ D
        R= 13
        LD@R 4
        ?A<C        MS          is it digit 0-3
        JNC-        BDD+A       no, go display address and data again
        GOSUB        APPHD
        JNC-        STP2        no hex digit, go display address and data again
*
* if key is accepted we return here
*
DP2D    ?NC XQ      NEXT2      add 2 prompts and wait for key
        JNC+        6          process backarrow key
*
        GOSUB        APPHD
        JNC-        DP2D       no hex digit, try with 2 prompts again
*
        JNC+        DP1D       skip backarrow processing
*
        READ D
        JNC-        STP2       remove data digit
                                display address and data again
*
DP1D    ?NC XQ      NEXT1      add 1 prompt and wait for key
        JNC+        6          process backarrow key
*
        GOSUB        APPHD
        JNC-        DP1D       no hex digit, try with 1 prompt again
*
        JNC+        3          skip backarrow processing
*
        READ D
        JNC-        DP2D       remove last digit
                                try with 2 prompts again
*

```



```

R= 13
LD@R 2
A=C      MS
GOSUB    RHEXD2
GOSUB    NULL
JNC-     STP2    key hold to long, display address + data again
?NC XQ   ENCP00
A<>C     A
RCR 3
A=C      X      save input data in "A(X)"
READ Q   get address to write to
FETCH
A<>C     X
A<>B     X      save original data word in "B(X)"
A=C      X      copy input data to "A(X)"
WRIT
FETCH
?A#C     X      is it written
JC+      NWBRAM  no, test for no write or bad mldl
DONDW1   READ Q   yes
C=C+1    M      point to next address
DONDW    A=C      M      save address in "A(M)"
GOTO     ANDW
*
NWBRAM   A<>B     X
?A#C     X      is the original data still the same
JC+      BMLDL   no, go say bad mldl
?NC XQ   CLLCDE
?NC XQ   MESSL
RMB      $08     "NO WRITE"
*
* put message in the display, wait for .3 seconds and ask for the address
* again
*
GOTO     ADLAG
*
RMWBSST  ?NC XQ   ENCP00
?NC XQ   OFSHFT  clear shift annunciator
READ Q
C=C-1    M      decrement address
JNC-     DONDW   go display address + data and wait for new data
*
* in case we are in rom 0 at word 0 and a BST is wanted, we would wrap to
* address FFFF. to avoid this we do a single step to warn the user
*
RMWSST   ?NC XQ   ENCP00
JNC-     DONDW1
*
* put message in display and reset partial key sequence and exit to the
* mainframe
*
BMLDL    ?NC XQ   CLLCDE
?NC XQ   MESSL
RMB      $08     "BAD MLDL"
?NC XQ   LEFTJ
SETF 8
?NC XQ   MSG105
?NC XQ   RSTSEQ
?NC XQ   STMSGF
?NC GO   NFRKB

```

```

*
*
*
RMWSTO  GOSUB      DPADR
        ?NC XQ     ENCP00
        READ Q
        FETCH      get data at current address
        C=C+1      M
        WRIT Q     save updated address
        RCR 10
        R= 13
        LD@R 2
        A=C        A
        ?NC XQ     ENLCD
        GOSUB      DPBYT  append data at this address to the display
        ?NC XQ     LEFTJ  view data PRINTDISPLAY
        LDI        $250   load timing constant
        CLRKEY
WAIKEY  ?KEY
        JNC+       NOKEY
        C=KEY      there is a key pressed
        RCR 3
        A=C        X
        A=0        XS    keycode to "A(X)"
        R= 0
        C=C+C      R    is it the ON key
        ?C GO      OFF   yes, go to deepsleep
        LDI        $087
        ?A#C       X    is it the R/S key
        JNC+       DOR/S yes, go process it
        RCR 11     no, get timing constant back
NOKEY   C=C-1      X    is time expired
        JNC-       WAIKEY no, keep waiting
        JNC-       RMWSTO yes, display next byte
*
DOR/S   ?NC XQ     RSTKB  wait until key is released
        LDI        3E8
        C=C+C      X
        C=C-1      X
        JNC-       1     wait a while
        GOTO      RMWBST do a bst and display address and data again
*
*
*
* Hex TO Decimal ( floating point )
*
* Entry : "A(X)" holds hex number ( smaller then 999 )
* Exit  : "A" , "C" holding floating point number
* Uses  : "A" , "B(X)" , "C"
*

```

```

HTOD      A=0      M
          A=0      MS      initialize "A[13-3]"
          LDI      $064
          C<>B      X      100d in "B(X)"
          C=0      A
D 100     ?A<B      X
          JC+      DO 10
          A=A-B      X      subtract 100d
          C=C+1      X      count the Hundreds in "A(X)"
          JNC-      D 100
DO 10     C<>B      A
          LDI      $00A
          C<>B      A      10d in "B(X)"
          RCR 13
D 10     ?A<B      X
          JC+      NORM
          A=A-B      X      subtract 10d
          C=C+1      X      count the Tens in "A(X)"
          JNC-      D 10

*
* normalize the number by setting up an exponent of 2 and shifting the
* mantissa to the left left one place and decrementing the exponent as long
* as the mantissa isn't in place
*
NORM      RCR 13      Units are left in "A(X)"
          A=A+C      X      add H,T and U to a dec. number
          LDI      $002      normalize number
          A<>C      A
          ?C#0      XS
          JC+      DONEHD
          RCR 13
          A=A-1      X
          ?C#0      XS
          JC+      DONEHD
          RCR 13
          A=A-1      X

*
* make the floating point number by combining the exponent and the mantissa
*
DONEHD    RCR 4      put mantissa in place
          C=C+A      X      add the exponent
          A=C      A      save floating point in "A" and "C"
          RTN

*
*
*
* COD
*
* Entry: alpha has the hex string that has to be converted
* Exit : NNN is in "B" and "C"
* Uses : "A", "B", "C"
*

```

| | | | |
|--|---------|--------|--|
| COD | B=0 | A | Initialize "B" for the new hex number |
| | SETF 9 | | Say we do the loop for the first time |
| | SLCT Q | | Select the loop counter |
| | READ M | | get first 7 digits to convert |
| | JNC+ | 3 | start converting loop |
| DO N | READ N | | get the next 7 digits |
| | CLRF 9 | | say we are finished after this loop |
| | A=C | A | save digits in "A" |
| | R= 7 | | set the loop counter for 7 digits |
| * | | | |
| * converting the ascii is done by testing if it lays in between \$30 and | | | |
| * \$39 and if so, subtracting \$30. if it isn't a number we test for ascii | | | |
| * between \$41 and \$46 and if so subtract \$37 | | | |
| * | | | |
| COD 1 | SLCT P | | |
| | R= 1 | | say we only compare the right digit |
| | ?A#0 | R< | add a zero to the NNN in "B" |
| | JNC+ | DO-9 | |
| | LDI | \$030 | load test constant for ascii < \$30 |
| | A=A-C | R< | |
| | JC+ | DATERR | we do have ascii < \$30 |
| | LDI | \$00A | |
| | ?A<C | R< | test if this is ascii for a digit 0-9 |
| | JC+ | DO-9 | we do have a digit 0-9 |
| | LDI | \$011 | |
| | ?A<C | R< | test if the ascii code is less then \$41 |
| | JC+ | DATERR | |
| | LDI | \$017 | test for ascii codes > \$46 |
| | ?A<C | R< | |
| | JNC+ | DATERR | |
| | LDI | \$007 | |
| | A=A-C | R< | make the hex value for a digit A-F |
| DO-9 | R= 0 | | |
| | B=A | R | add this digit to the NNN in "B" |
| | C<>B | A | |
| | RCR 1 | | |
| | C<>B | A | shift NNN one place left |
| | RSHFA | A | |
| | RSHFA | A | get next digit to "A[1-0]" |
| | SLCT Q | | |
| | R=R-1 | | |
| | ?R= 0 | | have we done 7 digits |
| | JNC- | COD 1 | no |
| | ?FSET 9 | | do we have to do 7 more digits |
| | JC- | DO N | yes |
| CODEND | C=B | A | copy the NNN to "C" |
| | RTN | | |
| DATERR | ?NC XQ | ERROR | |
| | FCB | \$022 | specify message 'DATA ERROR' |
| * | | | |
| * | | | |
| * | | | |
| * Decode subroutine | | | |
| * | | | |
| * Entry: "B" holds NNN to decode | | | |
| * flag 8 must be set and 9 must be cleared | | | |
| * Exit: Alpha register holds the decoded NNN | | | |
| * | | | |

```

DECOD  SLCT Q
      JNC+          3
DECOD2 WRIT N      save the last 7 digits
      CLRf 8      say we do the loop for the last time
      R= 7        set loop counter to 7
*
* decoding a hex digit to it's ascii representation is done by adding $30
* if the hex number is smaller then $0A and adding $37 if it is greater
* then $09
*
DECOD1 LSHFA      A
      LSHFA      A      make room for the next ascii coded digit
      C<>B      A
      RCR 13
      C<>B      A      get next digit to transform
      SLCT P
      R= 0          point to digit to transform
      A<>B      R      digit to "A[0]"
      ?A#0      R      if it is a leading zero, it must be suppressed
      JNC+          2
      SETF 9        say we have had all leading zero's
      ?FSET 9      do we have a leading zero
      JNC+          DIG=0 yes
      R= 1
      LDI          $00A
      ?A<C      R<      is it a number
      JC+          4      yes
      LDI          $007
      A=A+C      R<      make it a digit A-F minus $30
      LDI          $030
      A=A+C      R<      make it a ascii coded digit
DIG=0  A<>C      A
      A=C          A      copy ascii string to "C"
      SLCT Q
      R=R-1
      ?R= 0        are we finished yet
      JNC-          DECOD1 no, loop again
      ?FSET 8      have we done the loop twice
      JC-          DECOD2 no, save 7 digits in N and do the rest
      ?FSET 9      was the NNN zero
      JC+          3      no
      LDI          $030      yes, load only one zero
      WRIT M      save last 7 digits in M
      ?NC GO      XAVIEW display and return to mainframe
*
*
* MOVE
*
      NAM          MOVE
*
      GOSUB      COD      get the desired addresses to "C"
      A=0          A
      B=0          A      make sure we only leave the wanted addresses
      SLCT Q
      R= 6
      SLCT P
      R= 3          devine the address field with P and Q

```

```

*
* split the begin, end, and destination addresses
*
      RCR 1
      A<>C      PQ          EEEE to "A"
      RCR 4
      C<>B      PQ          BBBB to "B"
      RCR 6
      C<>B      PQ          DDDD to "B" BBBB to "C"
      ?A<B      M          test if EEEE<DDDD
      JC+              MOVE2
      A<>B      M          DDDD to "A" EEEE to "B"
      ?A<C      M          test if DDDD<BBBB
      JC+              MOVE3
*
* if the destination address lays in between the begin and end address of
* the block we want to move, we have to start copying at the begin of the
* source block and place it at the begin of the destination block because
* we overwrite the source block
*
      M=C              save BBBB in "M"
      A<>B      M
      C=A-C      M
      A<>B      M          number of bytes to copy to "C"
      A=A+C      M          end of destination block to "A"
      C=M              BBBB to "C"
      A<>B      M          end of dest to "B"
      A<>C      M          EEEE to "C" BBBB to "A"
      C=C+1      M          this is necessary to get all bytes copied
MOVE1  C=C-1      M          decrement source address
      FETCH              get source byte
      C<>B      M          get destination address
      WRIT              write byte to destination
      C=C-1      M          decrement destination address
      C<>B      M          get source address back to "C"
      ?A<C      M          are we finished yet
      JC-              MOVE1  no, loop again
      RTN
*
* when the destination block is not overwriting the source block we just
* copy from begin to end
*
MOVE3  A<>B      M          EEEE to "A" DDDD to "B"
MOVE2  FETCH              get source byte
      C=C+1      M          increment source address
      C<>B      M          get destination address
      WRIT              write byte to destination
      C=C+1      M          increment destination address
      C<>B      M          get source address back to "C"
      ?A<C      M          are we finished yet
      JNC-              MOVE2  no, loop again
      RTN
*
*
*
* Last ROM word
*
      NAM              LROM
*

```

```

        GOSUB      COD      get start address to "C[3-0]"
        A=C        M        get page digit to "A(M)"
        LSHFA      M
        LSHFA      M
        LSHFA      M        page digit to "A[6]"
        RCR 11     M        start address to "C[6-3]"
LROM1    C=C-1     M        point to next word to examine
        FETCH
        ?C#0       X
        JC+        LROM2    no, we have the first non zero word
        ?A<C       M        have we searched the whole rom yet
        JC-        LROM1    no, loop again
NONE     ?NC XQ    ERROR    specify message 'NONE'
        FCB        $031
LROM2    A=0       A        initialise for DECOD and put address and byte
        B=0       A        in alpha
        C<>B      A
        SETF 8
        CLRF 9
        WRIT 0
        WRIT P
        GOTO      DECOD

*
*
*
* LOCate by Alpha
*
        NAM        LOCA
*
        GOSUB      COD      get start address and data word
        R= 6
        A=0        A
        A=A-1      R<      create FFF in "A[5-3]"
        A=C        R        create the end address of this page in "A[6-3]"
        A=C        X        save the target word in "A(X)"
LOCA     C=C+1     M        point to next word to test
        FETCH
        ?A#C       X        is it the desired data word
DECODE   JNC-      LROM2    yes, send it to alpha and the display
        ?A<C       M        have we reached the end of the rom yet
        JNC-      LOCA     no, loop again

*
* set user flag 10 when we do not find the target word
*
        READ D
        RCR 11     get user flags
        C<>ST      user flag 10 to CPU flag 1
        SETF 1     set user flag 10
        C<>ST
        RCR 3
        WRIT D
        JNC-      NONE     restore the updated user flags
                                and say 'NONE'

*
*
*
* CODE
*
        NAM        COD
*

```

| | | | | |
|-------|-------------|----|--------|--|
| | GOSUB | | COD | get the NNN to "C" and "B" |
| | ?NC GO | | RCL | put the NNN on the user stack and take care of a stacklift |
| * | | | | |
| * | | | | |
| * | | | | |
| * | DECODE | | | |
| * | | | | |
| | NAM | | DECOD | |
| * | | | | |
| | READ X | | | |
| | JNC- | | DECODE | |
| * | | | | |
| * | | | | |
| * | CLear BLock | | | |
| * | | | | |
| | NAM | | CLBL | |
| * | | | | |
| | GOSUB | | COD | get the desired page or block to "C" |
| | RCR 12 | | | |
| | ?C#0 | M | | do we have to do the entire page |
| | JC+ | | CLBL2 | no, only a little block |
| | RCR 10 | | | get page address to "C[6]" |
| | R= 5 | | | |
| | A=C | A | | page to "A[6]" and clear the rest of "A" |
| | A=A-1 | R< | | create end address of this page |
| | A<>C | M | | EEEE to "C" BBBB to "A" |
| CLBL1 | WRIT | | | clear this word |
| | C=C-1 | M | | point to next word |
| | ?A<C | M | | are we finished yet |
| | JC- | | CLBL1 | no, loop again |
| | WRIT | | | clear the first word of the block |
| | RTN | | | |
| CLBL2 | RCR 3 | | | BBBB to "A(M)" |
| | A=0 | M | | |
| | A<>C | M | | |
| | RCR 10 | | | EEEE to "C(M)" |
| | JNC- | | CLBL1 | clear the block |
| * | | | | |
| * | | | | |
| * | | | | |
| * | Count ByTes | | | |
| * | | | | |
| | NAM | 10 | CBT | indicate an alpha prompt |
| * | | | | |
| | FCB | | \$000 | indicate not programmable |
| * | | | | |
| | GOSUB | | FPAL | |
| | SETHEx | | | |
| | RCR 8 | | | |
| | A=C | A | | address of end to "A[3-0]" |
| | CLRf 10 | | | say we are in ram |
| | N=C | | | save end address in "N" |
| | ?NC XQ | | CPGMHD | get start address to "A[3-0]" |
| | ?NC XQ | | PUTPCF | place PC to start of program |
| | C=N | | | |
| | ?NC XQ | | CALDSP | # bytes*2 in "A[3]" # regs in "A(X)" |
| | A<>C | R< | | |

| | | |
|--------|--------|--------------------------------|
| A=C | R< | |
| C=C+C | X | |
| C=C+C | X | |
| C=C+C | X | |
| A<>C | X | |
| A=A-C | X | compute registers * 7 |
| RCR 1 | | |
| C=0 | M | |
| RCR 2 | | |
| C=C+C | X | |
| C=C+C | X | |
| C=C+C | X | |
| RCR 1 | | right justify number of bytes |
| C=C+A | X | |
| C=C+1 | X | |
| C=C+1 | X | total program length to "C(X)" |
| A=0 | A | |
| A=C | X | |
| ?NC XQ | CLLCDE | |
| ?NC XQ | GENNUM | get number of bytes to display |
| LDI | \$020 | |
| WRIT E | | append a space |
| ?NC XQ | MESSL | |
| RMB | \$06 | "BYTES_" |
| ?NC XQ | LEFTJ | |
| LDI | \$020 | |
| WRIT D | | |
| ?NC XQ | ENCP00 | |
| ?NC XQ | STMSGF | |
| ?NC GO | NFRKB | |

*
*
*
*
*
*

Delete FAT entry

| | | |
|--------|------|--|
| NAM | DFAT | |
| GOSUB | COD | get entry address to "C[4-0]" |
| R= 1 | | set pointer for comparing digit 0 and 1 |
| C<>B | X | save AAA in "B(X)" |
| C=0 | X | |
| RCR 11 | | first address of this page to "C[6-3]" |
| A=C | A | save also in "A" |
| C=C+1 | M | |
| FETCH | | get number of functions on this page |
| C=C+C | X | double it to make number of bytes in FAT |
| A<>C | X | |
| RCR 3 | | |
| C=C+A | X | create address of last FAT word |
| RCR 11 | | |
| A<>C | M | end address to "A(M)" |
| C<>B | X | get AAA to "C(X)" |
| A=0 | X | |
| A<>C | XS | |
| RSHFA | X | |
| RSHFA | X | first digit of AAA in "A[0]" |
| RCR 6 | | get 0 to "C[1]" |
| A=C R | | complete first word of FAT in A"(X)" |
| RCR 8 | | |
| C<>B | X | second word of FAT in B"(X)" |

```

*
* in this loop we find the address of the first entry word of the entry
* to delete
*
DFAT1  C=C+1      M          point to word before an entry
      A<>B        X          these two instructions are necessary to assure
DFAT2  A<>B        X          proper operation in case we compare both words
      C=C+1      M          point to first FAT word
      ?A<C       M          have we had all entry's
      JC+                NO ENT yes, say there is not such an entry
      FETCH                get first FAT word
      ?A#C        R<        is it the same as our first FAT word
      JC-                DFAT1 no, do the next entry
      A<>B        X
      C=C+1      M
      FETCH                get second FAT word
      ?A#C        R<        is it the same as our second FAT word
      JC-                DFAT2 no, restore our FAT words and do the next entry
*
* after we have found the FAT entry to delete, we have to move the whole
* FAT down by two words and make sure that the last entry is replaced by
* the two necessary null words to say that this is the end of the FAT
*
      A=A+1      M
      A=A+1      M          point to the last null word after the FAT
      C=C-1      M          point to the first word of the undesired entry
      C<>B        M
      C=B        M          save destination address in "B(M)"
      C=C+1      M
DFAT3  C=C+1      M          increment source address
      FETCH                get source word
      C<>B        M
      WRIT                write it to destination address
      C=C+1      M          increment destination address
      C<>B        M          destination to "B(M)" source to "C(M)"
      ?A<C       M          are we finished yet
      JNC-                DFAT3 no, do next FAT word
      R= 5
      LD@R 0
      LD@R 0
      LD@R 1
      FETCH                get number of functions on this page
      C=C-1      X          decrement number of functions
      WRIT                save new number of functions
      RTN
NO ENT  GOSUB                INERR
      RMB                $08 "NO ENTRY"
DFERRE  GOTO                ERREX
*
*
*
* Append entry to FAT
*
*

```

```

* AFAT1
*
* Entry: "A" has the entry as follows in "C[5-0]" UOPAAA
*       "B(X)" has the digits AAA
* Exit : the entry is appended when there is enough room in the FAT, other
*       wise the message ENTRY>64 is displayed
* Uses : "A", "B(X)", "C"
*
*       NAM           AFAT
*
*       GOSUB         COD      get UOPAAA to "C" and "B"
*       A<>C           A
AFAT1   A<>C           A
*       A=C           M      save user code and offset in "A[5-4]"
*       C=0           X
*       RCR 11        M      address of first word to "C[6-3]"
*       C=C+1         M
*       FETCH         M      get number of functions
*       C=C+1         X      add one function
*       A<>C           X
*       LDI           $040    this is max number of functions on one page
*       A<>C           X
*       ?A<C          X      can we append another function
*       JC+           E>64    no, say 'ENTRY>64'
*       WRIT          M      increment the number of functions with one
*       C=C+C         X      this is address of first word of new function
*       A<>C           X
*       RCR 3         M
*       A<>C           X
*       RCR 11        M      "C[6-3]" has first word address
*       C=B           X      get AAA to "C(X)"
*       RSHFA         A      offset and user to "A[4-3]"
*       A=0           X
*       A<>C         XS      first digit AAA to "A[2]" second word to "C(X)"
*       RSHFA         A
*       RSHFA         A      first entry word in "A(X)"
*       A<>C           X
*       WRIT          M      append first entry word to FAT
*       C=C+1         M
*       A<>C           X
*       WRIT          M      append second entry word to FAT
*       C=0           X      make null word after FAT in "C(X)"
*       C=C+1         M
*       WRIT          M      first null word after FAT in RAM
*       C=C+1         M
*       WRIT          M      second null word
*       RTN
E>64   GOSUB         INERR
*       RMB           $08      "ENTRY>64"
*       JNC-         DFERRE
*
*
*
* --
*
*       NAM           --
*
*       RTN
*

```

```

*
*
*      ORG          $7FF
*
*
*
* MNEM
*
*      NAM          MNEM
*
*      C=0          A
*      RAMSLCT      select status registers
*      PRPHSLCT     and deselect all peripherals
*      READ L       get an instruction byte to "C[2-0]"
*
* when this is the first byte of an instruction there will happen nothing
* assumed is that reg L is empty when we start. when it is the second byte
* of an instruction the first byte is found in "C[2-0]". it is used to find
* out to which class the instruction belonged and what to do now
*
*      C<>ST        instr. class of first byte to flag 1 and 0
*      READ Y       read first byte of instruction
*
* here is tested if we have to do with a class 1 instruction or if we
* handle a LDI instruction
*
*      ?FSET 0      is it a class 1 instruction
*      JC+          STP 1 yes, handle this instruction type
*      ?FSET 1      is it a LDI
*      JC+          STP 2 it is a LDI, go handle it
*
* here is the class tested in case we do the first byte of an instruction
*
*      ST=C         get class to flag 1 and 0
*      ?FSET 0      is it class 1 or 3
*      JC+          STP 3 yes, go handle class 1 or 3
*      ?FSET 1      is it class 0 or 2
*      JNC+         STP 4 go handle class 0
*
* this part takes care of the arithmetic instructions
*
*      A=C          A
*      B=A          A      save instruction code and it's address in "B"
*      C=C+C        A
*      C=C+C        A
*      RSHFC        X      strip of the class bits
*      C=C+C        A
*      RSHFC        X      strip off the field bits and right justify
*      A=C          A      save number of instruction in "A(X)"
*      ?NC XQ       PCTOC  get current page address
*      R= 5
*      LD@R A
*      LD@R 0
*      LD@R 0          load start-1 of instruction table
*
* all the instructions in class two are numbered in sequence, determined
* by the first 5 bits of the instruction code. we pick the right point in
* the table by decrementing the number every time we have had an ascii
* string representing one instruction and exit when we have reached the
* right string
*

```

| | | | |
|--|---------|--------|--|
| | R= 1 | | set pointer for instruction counter |
| | CLRF 9 | | say we only are doing the instruction type |
| CLS 23 | A=A-1 | R< | test if we have reached our string number |
| | JC+ | | CLS 20 we have found our string |
| | C=C+1 | M | |
| | FETCH | | get next ascii character of a string |
| | C=0 | R< | |
| | C=C-1 | X | is this the end of a string |
| | JC- | 4 | no, pick next character of string |
| | JNC- | CLS 23 | yes, go test if we are ready yet |
| * stepping stones | | | |
| STP 4 | JNC+ | STP 5 | |
| STP 3 | JNC+ | STP 6 | |
| STP 2 | JNC+ | DOLDI | |
| STP 1 | JNC+ | STP 7 | |
| * we enter the fetch of the string we want to have in "A" with the address | | | |
| * of the first character-1 in "C(M)" | | | |
| CLS 20 | A=0 | A | initialize "A" for holding a string |
| | R= 1 | | look only for the ascii byte |
| CLS 21 | C=C+1 | M | |
| | FETCH | | get next character |
| | ST=C | | save it in "ST" |
| | LSHFA | A | |
| | LSHFA | A | |
| | A<>C | R< | add character to the string |
| | C=C-1 | X | have we done all characters |
| | JC- | CLS 21 | no, do next character |
| | A=A+1 | MS | type string as alpha data |
| | A<>C | A | and put string in "C" |
| | ?FSET 9 | | have we had the 'field' string yet |
| | JC+ | CLS 22 | yes, write field string to stack and exit |
| | SETF 9 | | no, say we do the field string |
| | M=C | | save instruction string in "M" |
| | C=B | A | get instruction code back |
| | C=C+C | X | |
| | C=C+C | X | |
| | RSHFC | X | strip off the class bytes |
| | C<>ST | | |
| | CLRF 3 | | |
| | C<>ST | | clear last bit of instruction type |
| | A=C | X | save field byte in "A(X)" |
| | ?NC XQ | PCTOC | |
| | R= 5 | | |
| | LD@R A | | |
| | LD@R B | | |
| | LD@R A | | load start address of field table |
| PARAM | R= 1 | | |
| | A=0 | R | leave the field byte only in "A[0]" |
| | JNC- | CLS 23 | and make the field string in "A" |
| CLS 22 | WRIT T | | write field string to reg T |
| | C=M | | |
| | A=C | A | instruction string to "A" |
| | LDI | \$020 | load ascii byte for a space |

```

CLS 24  R= 11
        ?A#0      R      do we have a character in byte 6
        JC+        CLS 25  yes, then we are finished
        LSHFA      R<
        LSHFA      R<      no, make place for a space at the right end
        R= 1
        A=C        R<      add space to string
        JNC-        CLS 24  and test if we are done
CLS 25  A<>C      A
        WRIT Z      save field string in reg Z
        C=0        A
        WRIT L      say we have first byte of instruction
        RTN

*
* stepping stones
*
STP 5   JNC+      STP 8
STP 6   JNC+      STP 9
STP 7   JNC+      DOCLS1
*
*
* this part converts the hex constant of the LDI instruction to its decimal
* value in "A[3-1]"
*
DOLDI   R= 2      set loop counter for conversion
        A=0      A      initialize "A" for conversion
LDI2    C=C-1     R      have we done this digit
        JC+      LDI1    yes, test if we are ready
        SETDEC
        A=A+1     A      add converted hex value to "A"
HMNDIG  SETHEX
        JNC-      LDI2
LDI1    ?R= 0     are we finished yet
        JC+      CONRDY
        R=R-1
        C<>B      A      no, decrement loop counter and multiply by 16
                        save hex number in "B(X)"
*
* the conversion is done by counting how many times we have a factor of 16
* in the hex number. this is done in decimal. after we have found this
* number, we multiply it by 16 and add the rest of the hex number to the
* decimal number. this leaves the converted result in "A"
*
        A<>C      A
        A=C        A      copy decimal string to "C"
        LSHFA      A      multiply "A" by 10
        SETDEC
        C=C+C      A
        A=A+C      A
        A=A+C      A
        A=A+C      A      "A" is multiplied by 16
        C<>B      A      get hex number back
        JNC-        HMNDIG  set mode to hex and do next digit
CONRDY  A<>C      A
*
* at this point "C[3-0]" contains the converted hex string right justified
* we only have to make it an ascii coded decimal string. this is done by
* inserting the digit "3" before every decimal digit
*

```

| | | | |
|---|---------|--------|--|
| | RCR 11 | | initialize string for ascii conversion |
| | R= 3 | | set loop counter |
| LDI3 | RSHFC | R< | |
| | LD@R 3 | | convert one digit to ascii |
| | R=R+1 | | |
| | R=R+1 | | increment loop counter |
| | ?R= 6 | | are we done yet |
| | JNC- | LDI3 | no, loop again |
| | R= 7 | | |
| | LD@R 3 | | do the last character |
| | C=C+1 | MS | type string as alpha data |
| | WRIT T | | and put it in reg T |
| | C=0 | M | |
| | R= 11 | | |
| | LD@R 4 | | |
| | LD@R 3 | | |
| | LD@R 4 | | |
| | LD@R F | | |
| | LD@R 4 | | |
| | LD@R E | | |
| | LD@R 2 | | |
| | LD@R 0 | | |
| | LD@R 2 | | |
| | LDI | \$020 | load the 'CON___' string |
| | WRIT Z | | and save it in reg Z |
| | C=0 | A | |
| | WRIT L | | say we are doing the first byte |
| | RTN | | |
| * * stepping stones * | | | |
| STP | 8 | JNC+ | STP 10 |
| STP | 9 | JNC+ | STP 11 |
| * * | | | |
| * here is the class 1 instruction type field handled and the jump | | | |
| * address is done | | | |
| * | | | |
| DOCLS1 | ST=C | | get the second byte into "ST" for testing type |
| | A=C | A | save first byte in "A(X)" |
| | R= 13 | | |
| | LD@R 1 | | |
| | LD@R 0 | | load alpha identifier |
| | LD@R 4 | | |
| | LD@R 7 | | |
| | LD@R 4 | | |
| | LD@R F | | first two letters are always 'GO' |
| | ?FSET 1 | | is it a jump or an xeq |
| | JNC+ | DOXEQ | it is an xeq |
| | LD@R 4 | | |
| | LD@R C | | load a 'L' for the 'GOL' |
| | ?FSET 0 | | is it a jump on carry |
| | JC+ | GOLC | yes, handle this one |
| | LD@R 4 | | |
| | LD@R F | | |
| | LD@R 4 | | |
| | LD@R E | | |
| | LD@R 4 | | |
| | LD@R 7 | | finish the instruction to 'GOLONG' |
| | JNC+ | JMPTOZ | |

| | | | |
|--|---------|--------|---|
| DOXEQ | LD@R 5 | | |
| | LD@R 3 | | make the string 'GOS' |
| | ?FSET 0 | | is it an xeq on carry |
| | JNC+ | GOSUB | no, handle this one |
| GOLC | LD@R 4 | | |
| | LD@R 3 | | |
| | LD@R 2 | | |
| | LD@R 0 | | make instruction 'GOSC_' or 'GOLC_' |
| | JNC+ | ADDSP | add a space for 6 characters total |
| | LD@R 5 | | |
| | LD@R 5 | | |
| | LD@R 4 | | |
| | LD@R 2 | | make string 'GOSUB' |
| | LD@R 2 | | |
| ADDSP | LD@R 0 | | add a space to the string to get 6 characters |
| JMPTOZ | WRIT Z | | save string in reg Z |
| | A<>C | X | get the two msb digits of address to "C[2-0]" |
| | C=C+C | A | |
| | C=C+C | A | |
| | RCR 12 | | shift them into "C[2-1]" |
| | A<>C | A | save them in "A(M)" |
| | READ L | | get the lsb digits |
| | A<>C | X | |
| | A<>C | A | |
| | C=C+C | X | |
| | C=C+C | X | left justify them in "C(X)" |
| | RSHFC | A | right justify address in "C" |
| | CLRF 9 | | say we are doing a class 1 type instruction |
| | JNC+ | DIST | skip the stepping stones |
| * * stepping stones * | | | |
| STP 10 | JNC+ | STP 12 | |
| STP 11 | JNC+ | CLS1FB | |
| * * * in this part we handle the addresses of the jump and gosub instructions * and we handle the distance part of the short jumps. when we have a short * jump we also add a asterisk and the direction * first thing we have to do for this is to translate the hex representation * of the address or the distance to the ascii coded hex form * | | | |
| DIST | R= 3 | | set loop counter to the starting position |
| | A<>C | R< | |
| | C=0 | A | |
| | A<>C | R< | |
| | RCR 11 | | |
| | A<>C | A | |
| | C=0 | A | |
| | SETDEC | | |
| | C=-C-1 | A | |
| | SETHex | | |
| | A<>C | A | get all 9's to "A" and address to "C" |
| NDIG | RSHFC | R< | shift digit one right |
| | LD@R 3 | | and make it ascii '0'-'9' |
| | ?A<C | R | was it a number |
| | JNC+ | NUMBER | yes, then we are ready |
| | C=A-C | R | |
| | C=-C | R | no, convert A-F to 1-6 |
| | R=R+1 | | |
| | LD@R 4 | | and make it ascii 'A'-'F' |

| | | | | |
|--|---------|--------|--|--|
| NUMBER | R=R+1 | | | increment loop counter by one |
| | R=R+1 | | | are we finished yet |
| | ?R= 6 | | | no, do next digit |
| | JNC- | NDIG | | |
| | R= 7 | | | |
| | LD@R 3 | | | |
| | ?A<C | R | | |
| | JNC+ | NUMB | | |
| | C=A-C | R | | |
| | C=-C | R | | |
| | R= 7 | | | |
| | LD@R 4 | | | |
| | ?FSET 9 | | | is it the address of a xeq or jump |
| | JNC+ | DONE | | yes, exit |
| | R= 7 | | | no, we have to add an asterisk and direction |
| | LD@R 2 | | | |
| | LD@R A | | | load the asterisk |
| | LD@R 2 | | | |
| | LD@R B | | | load direction forward |
| | ?FSET 8 | | | is it backwards |
| | JC+ | FWRD | | no, do forward |
| | R= 4 | | | |
| | LD@R D | | | load direction backwards '-' |
| FWRD | C=C+1 | MS | | type string as alpha data |
| | WRIT T | | | save it in reg T |
| | C=0 | A | | |
| | WRIT L | | | say we are doing the first byte |
| | RTN | | | |
| * | | | | |
| * stepping stones | | | | |
| * | | | | |
| STP 12 | JNC+ | CLS 0 | | |
| DIST1 | JNC- | DIST | | |
| * | | | | |
| * | | | | |
| * in this part we handle the class 1 instructions in case we have them for | | | | |
| * the first byte | | | | |
| * also are the class three instructions handled here | | | | |
| * | | | | |
| CLS1FB | ?FSET 1 | | | is it class 1 or 3 |
| | JNC+ | BYTE11 | | we have class 1 |
| | SETF 9 | | | say we do class 3 distance for converting part |
| | A<>C | A | | save instruction code in "A(X)" |
| | R= 13 | | | |
| | LD@R 1 | | | |
| | LD@R 0 | | | load alpha identifier |
| | LD@R 4 | | | |
| | LD@R 7 | | | |
| | LD@R 4 | | | |
| | LD@R F | | | make string 'GO' |
| | ?FSET 2 | | | is it a gonc |
| | JC+ | GOC | | no, do goc |
| | LD@R 4 | | | |
| | LD@R E | | | |
| | LD@R 4 | | | |
| | LD@R 3 | | | make string 'GONC' |
| | JNC+ | ADDSPA | | |
| GOC | LD@R 4 | | | |
| | LD@R 3 | | | |
| | LD@R 2 | | | |
| | LD@R 0 | | | make string 'GOC_' |

```

ADDSPA  LD@R 2
        LDI          $020      add two spaces to string
        WRIT Z        save string in reg Z
        A<>C          A        get instruction code back
        SETF 8        say we do a forward jump
        C=C+C          A
        RSHFC          X        strip of class and c/nc bits
        ST=C          direction bit to "ST 6"
        ?FSET 6        is it a forward jump
        JNC-          DIST1    yes, convert distance and say forward
*
* when we have a jump backwards we have to convert the coding, because the
* distance is saved in the inverted format
*
        SETF 7        make sure that the first bit becomes zero
        C=ST
        C=-C          X        invert the bit pattern
        CLRF 8        say we jump backwards
        JNC-          DIST1    convert distance and say backward
*
* here is the handling of class 1 done in case we have to do the first byte
* of the instruction
*
BYTE11  WRIT L        save instruction code in reg L for second pass
        C=0          A
        C=C+1        MS
        WRIT Z
        WRIT T        write a zero string to reg Z and T
        RTN
*
* here the handling of the class 0 instructions is done
*
CLS 0   C=C+C          X
        C=C+C          X
        RSHFC          X        strip off the class bits and right justify
        ST=C          instr. to "ST" for testing
        A=C            X        and save in "A(X)"
        CLRF 8
        CLRF 9        say we are handling row 6
        LDI          $006
        R= 0
        ?A#C          R        is it row 6
        JNC+          DOROW6    yes, go handle it
        LDI          $008
        ?A#C          R        is it row 8
        JNC+          DOROW8    yes, go handle it
        LDI          $00C
        ?A#C          R        is it row C
        JC+          STP14      no, we have not to do with row 6,8 or C
        SETF 9        say we do row C
        JNC+          DOROW6
DOROW8  SETF 8        say we do row 8
DOROW6  A<>B          X        save instr byte in "B(X)"
        C=B            X        copy instr. byte to "C(X)"
        RSHFC          X        strip of the row byte
        R= 1
        ?FSET 8        do we have to do row 8
        JNC+          2
        C=C+1          R        yes, add 16 to instr. byte
        ?FSET 9        are we doing row C

```

| | | | |
|------------------|--------|--------|--|
| | JNC+ | 3 | |
| | C=C+1 | R | |
| | C=C+1 | R | yes, add 32 to instr. byte |
| | A=C | X | save instr. byte in "A(X)" |
| | ?NC XQ | PCTOC | |
| | R= 5 | | |
| | LD@R A | | |
| | LD@R 9 | | |
| | LD@R 9 | | load start address of instr. of row 6,8,C |
| | R= 1 | | |
| | A=A-1 | X | have we found the wanted instruction |
| | JC+ | 7 | yes, get the string to "A" |
| | C=C+1 | M | |
| | FETCH | | |
| | C=0 | R< | |
| | C=C-1 | X | is this the end of an instruction |
| | JC- | 4 | no, try next character of instruction |
| | JNC- | 7 | yes, test if we are finished |
| | A=0 | A | initialize "A" for string |
| | C=C+1 | M | |
| | FETCH | | get a character |
| | LSHFA | A | |
| | LSHFA | A | |
| | A<>C | R< | append character to string |
| | C=C-1 | X | are we finished yet |
| | JC- | 6 | no, do next character |
| | A=A+1 | MS | type it as an alpha string |
| | LDI | \$020 | load ascii byte for a space |
| | R= 11 | | |
| | ?A#0 | R | is the string left justified |
| | JC+ | 6 | yes, exit |
| | LSHFA | R< | |
| | LSHFA | R< | |
| | A=A+C | X | no, shift one byte left and append a space |
| | JNC- | 5 | and try again |
| * | | | |
| * stepping stone | | | |
| * | | | |
| STP14 | JNC+ | DOCLS0 | |
| * | | | |
| | A<>C | A | |
| | WRIT Z | | save instruction string in reg Z |
| | A<>B | A | get instruction code into "A(X)" |
| | LDI | \$0DC | |
| | ?A#C | X | isi it a C=C or A |
| | JC+ | CANDA | no, test for C=C and A |
| | C=0 | A | |
| | R= 7 | | |
| | LD@R 4 | | |
| | LD@R F | | |
| | LD@R 5 | | |
| | LD@R 2 | | load 'OR' |
| APP_A | LD@R 2 | | |
| | LD@R 0 | | |
| | LD@R 4 | | |
| | LD@R 1 | | load '_A' |

| | | | | |
|--|--------|----|--------|--|
| EX1 | C=C+1 | MS | | type as alpha |
| | WRIT T | | | and save in reg T |
| | LDI | | \$04C | |
| | ?A#C | X | | is it a LDI |
| | JNC+ | | LDI1 | |
| C=0 | C=0 | A | | no, normal instruction go clear reg T and L |
| | A | | | |
| | WRIT L | | | |
| | RTN | | | |
| CANDA | LDI | | \$0EC | |
| | ?A#C | X | | |
| | JC+ | | A | |
| | C=0 | A | | |
| | R= 9 | | | |
| | LD@R 4 | | | |
| | LD@R 1 | | | |
| | LD@R 4 | | | |
| | LD@R E | | | |
| | LD@R 4 | | | |
| | LD@R 4 | | | load 'AND' |
| | JNC- | | APP_A | append '_A' and exit |
| | C=0 | A | | |
| | JNC- | | EX1 | |
| LDI1 | LDI | | \$002 | |
| | WRIT L | | | tell MNEM for second byte that we have a LDI |
| | RTN | | | |
| * | | | | |
| * here are the normal instructions from class 0 handled. these are the | | | | |
| * instructions that do have a parameter with them | | | | |
| * | | | | |
| DOCLS0 | ?NC XQ | | PCTOC | |
| | R= 5 | | | |
| | LD@R B | | | |
| | LD@R 8 | | | |
| | LD@R 5 | | | load start address of first string |
| FSTR | B=A | X | | save instr. code in "B(X)" |
| CLS 03 | R= 0 | | | use only the instruction number and not param. |
| | A=A-1 | R | | have we reached our instruction |
| | JC+ | | CLS 01 | yes |
| CLS 02 | C=C+1 | M | | |
| | FETCH | | | get next character |
| | R= 2 | | | |
| | ?C#0 | R | | is it the end of the string |
| | JNC- | | CLS 02 | no, try next character |
| | JNC- | | CLS 03 | yes, go test if we are finished |
| CLS 01 | R= 1 | | | set pointer to fetch and build instr. string |
| | A=0 | A | | |
| CLS 04 | C=C+1 | M | | |
| | FETCH | | | get a character of our string |
| | LSHFA | A | | |
| | LSHFA | A | | |
| | A<>C | R< | | append it to the string |
| | C=C-1 | X | | have we done all characters |
| | JC- | | CLS 04 | no, do next character |
| | A=A+1 | MS | | type string as alpha data |
| | A<>C | A | | |
| | M=C | | | and save it in "M" |
| | JNC+ | | B | |
| | LD@R 8 | | | |
| | LD@R 4 | | | |
| | LD@R C | | | |

| | | | | |
|------------------|---------|--------|--|---|
| | PUSH | | | |
| | C<>B | A | | |
| | RTN | | | |
| * | | | | |
| * stepping stone | | | | |
| * | | | | |
| STPP20 | JNC- | 6 | | |
| * | | | | |
| | WRIT Z | | | put string in reg Z |
| | ?FSET 9 | | | are we ready |
| | JNC+ | 5 | | no, continue |
| | C=0 | A | | |
| | C=C+1 | MS | | yes, clear reg T |
| | WRIT T | | | |
| | RTN | | | |
| * | | | | |
| | A<>B | A | | get instruction code back |
| | R= 1 | | | |
| | LDI | \$0F7 | | |
| | ?A#C | R | | is it a instr. from colum 15 |
| | JC+ | CLS 09 | | no, do parameter tables |
| | ?FSET 3 | | | is it row 8-F |
| | JC+ | CLS 09 | | yes, do parameter tables |
| | ?FSET 2 | | | is it row 0-3 |
| | JC+ | TST57 | | no, test for row 5 or 7 |
| | R= 0 | | | |
| | ?A#0 | R | | is it the NOP from row 0 |
| | JC+ | DOEXCP | | no, then it is the an exception |
| TST57 | ?FSET 0 | | | is it row 5 or 7 |
| | JNC+ | CLS 09 | | no, do parameter tables |
| DOEXCP | R= 1 | | | |
| | A=0 | R | | leave instr. byte only |
| | ?NC XQ | PCTOC | | |
| | R=5 | | | |
| | LD@R B | | | |
| | LD@R B | | | |
| | LD@R 6 | | | load address of exception row D table |
| | SETF 9 | | | say we don't have to do a parameter anymore |
| | JNC- | FSTR | | go get instruction string |
| * | | | | |
| CLS 09 | ?NC XQ | PCTOC | | |
| | R= 5 | | | |
| | LD@R B | | | |
| | LD@R D | | | |
| | LD@R 7 | | | load address of normal parameter table |
| | R= 0 | | | |
| | LDI | \$00A | | |
| | ?A#C | R | | do we have to do the parameters of REGN=C |
| | JC+ | 4 | | no |
| TREG | R= 4 | | | load table address of C=REGN and REGN=C |
| | LD@R F | | | |
| | JNC+ | REGN= | | finish address of table |
| | LDI | \$00E | | |
| | ?A#C | R | | do we have to do the parameters of C=REGN |
| | JNC- | TREG | | yes, load address of table |
| | LDI | \$004 | | |
| | ?A#C | R | | do we have to do the LD@R parameters |
| | JC+ | NPARAM | | no, do the normal parameters |

```

LD@R      R= 4
          LD@R E
REGN=     LD@R D      load table address of LD@R parameters
NPARAM    SETF 9      say we don't have to do a parameter anymore
          RSHFA      A      field byte to "A[0]"
          C<>B      A
          C=B        A      copy table address to "B(M)" page digit to "C"
          R= 5
          JNC-        STPP20 go do PARAM
*
*
*
          ORG          $A00
*
          ASCTBL
*
          RMB          $143
*
*
* the following ascii tables belong to the MNEM program. they have to start
* at the above specified address. the last character of every entry in the
* tables has bit 9 of the romword set. e.g. the entry A=0 is coded as
* follows A = 041
*          = = 03D
*          0 = 230
* the @ character means that there is just an empty entry in the table. it
* is coded with 200
*
CLASS2    A=0          these are the mnemonics for the class 2
          B=0          instructions
          C=0
          AB EX        the field specifier is picked from the next
          B=A          table
          AC EX
          C=B
          BC EX
          A=C
          A=A+B
          A=A+C
          A=A+1
          A=A-B
          A=A-1
          A=A-C
          C=C+C
          C=A+C
          C=C+1
          C=A-C
          C=C-1
          C=-C
          C=-C-1
          ?B#0
          ?C#0
          ?A<C
          ?A<B
          ?A#0
          ?A#C
          A SR
          B SR
          C SR
          A SL

```


| | | |
|--------|--------|---|
| CLSSOB | S=0 | this is the table of the regular class 0 |
| | S=1 | instructions. |
| | ?S=1 | these instructions can have 3 different |
| | LC | parameters. |
| | ?PT= | these parameters are specified in 3 tables |
| | @ | empty entry |
| | PT= | |
| | @ | empty entry |
| | SELP | |
| | REGN=C | |
| | ?F=1 | |
| | @ | empty entry |
| | C=REGN | |
| | RCR | |
| CLSSOC | @ | empty entry |
| | CLR ST | this is part of the class 0 special instruction |
| | RST KB | set. |
| | CHK KB | |
| | @ | empty entry |
| | DEC PT | |
| | @ | empty entry |
| | INC PT | |
| CLSSOD | 3 | class 0 parameter table for the following |
| | 4 | instructions |
| | 5 | NOP S=0 S=1 ?S=1 RCR ?F=1 |
| | 10 | |
| | 8 | SELP ?PT= PT= |
| | 6 | |
| | 11 | |
| | 14 | |
| | 2 | |
| | 9 | |
| | 7 | |
| | 13 | |
| | 1 | |
| | 12 | |
| | 0 | |
| | 15 | |
| CLSSOE | 0 | this is the table with the parameters for |
| | 1 | LC instruction |
| | 2 | |
| | 3 | |
| | 4 | |
| | 5 | |
| | 6 | |
| | 7 | |
| | 8 | |
| | 9 | |
| | A | |
| | B | |
| | C | |
| | D | |
| | E | |
| | F | |

| | | | |
|--------|-------|--|---------------------------------------|
| CLSSOF | 0(T) | | this is the parameter table for the |
| | 1(Z) | | instructions C=REGN and REGN=C |
| | 2(Y) | | |
| | 3(X) | | |
| | 4(L) | | |
| | 5(M) | | |
| | 6(N) | | |
| | 7(O) | | |
| | 8(P) | | |
| | 9(Q) | | |
| | 10(') | | the ' stands for the lazy T, code 07F |
| | 11(a) | | |
| | 12(b) | | |
| | 13(c) | | |
| | 14(d) | | |
| | 15(e) | | |

*
*
*
* ROM>REG
*

| | | | |
|--|-----|---------|--|
| | NAM | ROM>REG | |
|--|-----|---------|--|

*
* start initialization ROM>REG
*

| | | | |
|---------|----|--------|---|
| GOSUB | | LREG | get address of last existenet register |
| C=0 | X | | |
| RAMSLCT | | | |
| READ X | | | get start reg from X |
| ?NC XQ | | BCDBIN | |
| A=C | X | | save start reg number in "A(X)" |
| READ C | | | |
| RCR 3 | | | address of reg 00 to "C(X)" |
| C=C+A | X | | |
| M=C | | | address of start reg SSS to "M" |
| A=C | X | | and to "A(X)" |
| A<>B | X | | last reg to "A(X)" SSS to "B(X)" |
| C=A-C | X | | number of available registers to "C(X)" |
| A=C | X | | |
| C=C+C | X | | *2 |
| C=C+C | X | | *4 |
| A=A+C | X | | *5 |
| A=A-1 | X | | number of rom words that can be stored |
| B=A | X | | save max rom words in "B(X)" |
| READ Y | | | get BBBBEEEE |
| A=0 | A | | |
| R= 3 | | | |
| A<>C | R< | | EEEE to "A[3-0]" |
| RCR 4 | | | BBBB to "C[3-0]" |
| N=C | | | save BBBB in "N" |
| C=A-C | R< | | number of words to store NNN |
| LSHFA | A | | |
| LSHFA | A | | |
| LSHFA | A | | EEEE to "A(M)" |
| A=C | X | | NNN to "A(X)" |
| A<>B | X | | NNN to "B(X)" available to "A(X)" |
| ?A<B | X | | is there enough room to store NNN |
| ?C GO | | ERRNE | no, say 'NONEXISTENT' |
| C=M | | | |

```

A=C      X      SSS to "A(X)"
C=N
RCR 11
C<>B     M      BBBB to "B(M)"
C=B      A      x000000BBBBNNN to "C"
C=0      MS     clear the x "C" = 0000000BBBBNNN
RCR 7
A<>C     X      "C" = BBBBNNN0000000
N=C      X      "C" = BBBBNNN0000SSS
                        save header in "N"
*
* start main loop
*
* the main loop is entered with the CPU registers as follows
* "A(M)" = EEEE this is start address in rom
* "B(M)" = BBBB this is end address in rom
* "M(X)" = SSS this is the first MM register to use
* "N" = BBBBNNN0000SSS this is the header
*
MAINLP A<>B     A      EEEE to "B(M)" BBBB to "A(M)"
      R= 5      set loop counter to 5
ONEREG A<>C     A      address to "C" and string to "A"
      LSHFA     A
      LSHFA     A
*
* this makes place for the next rom word and shifts a zero into the
* alpha indicator digit, so we are sure that an increment of this digit
* returns the alpha indicator
*
      FETCH
A=C      X      put the rom word into the string
A<>C     A      string to "C" for bit shifting
C=C+C    X
C=C+C    X      push word up against rest of string
C=C+C    A
C=C+C    A      shift string left to even digit
A=A+1    M      increment address
R=R-1
?R= 0
JNC-      ONEREG no, do another rom word in this reg
RSHFC     A      right justify string in "C"
C=C+1    MS     set alpha indicator digit
C<>M
C=C+1     X      save string in "M"
RAMSLCT
C<>M      save reg address and get back string
WRITDAT
A<>B     A
?A<B     M      did we do all the rom words we had to do
JNC-      MAINLP no, continue
*
* termination of ROM>REG
*

```

```

C=M
A=C      X      last used register address to "A(X)"
C=N      get the header
RAMSLCT  and select the header register
C=A-C    X
C=C+1    X      number of used registers to "C(X)"
RCR 7     header is 0000RRRBBBBNNN   RRR is used regs
C=C+1    MS     make it alpha data
WRITDAT   save header in header register
C=0      X
RAMSLCT
READ C
RCR 3     reg 00 to "C(X)"
A=A-C    X      reg number of last used register
GOSUB     HTOD   make it floating point
WRIT L    and put it in L for the user
RTN

*
*
*
* REG>ROM
*
      NAM      REG>ROM

      READ Y
      N=C      start address to "N"
      READ X
      ?NC XQ   BCDBIN
      A=C      X      start register number to "A(X)"
      READ C
      RCR 3
      C=C+A    X      address of header register
      RAMSLCT
      M=C      save SSS
      READDAT  get the header register 100000AAAANNN
      RCR 3
      A=C      A      start address of file to "A[3-0]"
      C=N      get the wanted start address
      RCR 1
      C=0      M      leave only the start address
      ?C#0     X      is it a complete address
      JC+      8      yes, go put it in place
      RCR 10   put page number in "C[3]"
      A<>C     X      get address in this page
      R=3
      ?C#0     R      was there a page address given
      JC+      4      yes, put address in place
      A<>C     R      no, write it back at the original place
      JNC+     2
      RCR 13
      RCR 11
      A=C      A      put address in "A(M)"
      READDAT  get the header back again
      C=0      M
      RCR 11   leave only NNN in "C(M)"
      C=C+A    M
      C<>B     A      save end address in "B(M)"

*
* start of main loop
*

```

```

LOOP1  C=M
        C=C+1      X      point to next register
        RAMSLCT      select it
        M=C          save new register address
        READDAT      get 5 rom words
        RCR 13
        C=C+C        A
        C=C+C        A      left justify string in "C"
        R= 5          set LOOP2 counter for 5 words
LOOP2  RCR 12          shift 2 digits into "C(X)"
        C=C+C        A      shift one bit left
        JNC+          2      have we shift a one bit of at the left
        C=C+1        A      yes, append this bit at the right
        C=C+C        A
        JNC+          2
        C=C+1        A      do the second bit to
        A=C           X      rom word to "A(X)"
        A<>C         A
        WRIT          put the rom word in the ram
        A<>C         A
        ?A<B         M      have we done all the words
        ?NC RTN       yes, we are finished
        A=A+1        M      no, increment address
        R=R-1
        ?R=0          have we done the complete register
        JNC-          LOOP2 no, do the rest of this register
        JNC-          LOOP1 yes, do the next register

```

```
*
*
*
* these are a few subroutines that belong to the MMTORAM program
*
```

```

*
* this part takes care of the C handling
*
DO C      R= 0          set pointer to zero
          ?A<C          R      exceptions in C row
          ?NC RTN       return if exception
          C=M
          C=C+1         M      address of label to "C(M)"
          RCR 3         right justify address for AFAT1
          B<>C          A
          RCR 11        current PC to "B[6-3]" and AAA to "B(X)"
          C=B           X
          B<>C          A
          R= 5
          LD@R 2
          LD@R 0         say we have user code and offset is zero
          A<>C          A

*
* "B" has the current PC in [6-3] for the main loop of MMTORAM and the last
* three digits of the address in "B(X)" for AFAT1
* "A[5-0]" has UOPAAA for AFAT1
*
          GOSUB          AFAT1    append FAT entry and error if entrys > 64
          B<>C          A
          RCR 3
          B<>C          A      put PC back into "B[3-0]"
          RTN

*
*
*
* this part takes care of the B-E bytes ( 2 and 3 bytes GTO and XEQ and the
* alpha labels the last have to be appended to the FAT
*
DO B-E    LDI           $0CE      set up DE, C, B compare values
          ?A<C          R      is it B
          JNC+          NO B      no, continue
          GOTO          DO B      yes, handle B
NO B      ?A#C          R      is it C
          JNC-          DO C      yes, handle C

*
* handling the D bytes consist of converting the register length into a
* byte count, add the normal byte count to it and to invert the direction
* bit
*
          A<>C          X
          ST=C          save byte in "ST" for bit testing
          R= 3
          A<>B          R<
          ?NC XQ        NXBYTA    get PC back to "A[3-0]"
          A<>B          R<      get second byte of this instr.
          R= 2          save PC in "B[3-0]"
          C=0           R
          ?FSET 0       set up for conversion
          JNC+          MLTREG    is 1xx bit set
          C=C+1         R      yes, add 256 registers to the count
          CLR@ 0         leave byte count only in "ST"

*
* multiply the register count by 7 and add the byte count to it
*

```

```

MLTREG  A=C          X
        C=C+C        X
        C=C+C        X
        C=C+C        X
        A<>C         X
        A=A-C        X      regs * 7 in "A(X)" this is byte count of regs
        C=ST
        LD@R 0
        LD@R 0        get the byte count
        C=C+C        X
        C=C+C        X
        C=C+C        X
        C=0          M
        RCR 1         shift byte count three bits right
        A=A+C        X      total byte count of label in "A(X)"
        R= 2
        C=G          get identifier back to "C[3]"
        A<>C         X      add link bytes
        G=C          save first byte in "G"
        ST=C          and second in "ST"
        C=M
        C=C+1        M      get storage address
        LDI          $100    set first byte of instr. flag
        R= 0
        C=G          get first byte
        WRIT          and put it in ram
        CLRF 8        say we have had the first byte
        C=C+1        M      make storage address of second byte
        M=C          and save it
        C=ST          get second byte
        C=0          XS     clear first byte flag
        WRIT          and save it

*
* last thing to do here is to get the last byte and invert the direction
* bit
*
        R= 3
        A<>B        R<      get PC to "A[3-0]"
        ?NC XQ      NXBYTA  get the last byte
        A<>B        R<      save PC in "B[3-0]"
        ST=C          put byte in "ST" for inverting direction bit

*
* invert the direction bit
*
        ?FSET 7
        JC+          DIR1
        SETF 7
        JNC+         DIR2
DIR1     CLRF 7
DIR2     R= 0
RMLLOOP  C=ST
        G=C          save byte in "G"
        RTN          return to the main loop

*
*
* this part takes care of the handling of the B bytes
*

```

```

DO B      C=M
          C=C+1      M      get storage address
          M=C          save it
          R= 2
          LD@R 1      set first byte flag
          R= 0
          C=G          get first byte
          WRIT        put it in ram

*
* test direction bit of distance byte and compute the distance in bytes
*
          CLRf 8          use this flag for copy of direction bit
          R= 3
          A<>B      R<      get PC
          ?NC XQ      NXBYTA get the distance byte
          A<>B      R<      save PC
          ST=C        put distance byte in "ST" for testing direction
          ?FSET 7      is direction bit set
          JNC+        3
          SETf 8        yes, copy it into flag 8
          CLRf 7        clear distance byte for calculation of distance
          C=ST
          C=0      XS
          R= 1
          G=C          save bytes in "G"

*
* multiply register length by 7
*
          LD@R 0
          A=C      X
          C=C+C      X
          C=C+C      X
          C=C+C      X
          A<>C      X
          A=A-C      X      REG * 7 = bytes in "A[1-0]"
          C=0      X
          C=G          get bytelength to "C[0]"
          C=C+A      X      total byte length to "C[1-0]"
          ST=C
          ?C#0      X      is the distance to far
          JNC+        4      yes, leave direction bit zero
          ?FSET 8      was direction bit set
          JC+        2      yes, invert it
          SETf 7
          CLRf 8
          JNC-      RMLLOOP      say we have had the first byte
                                save byte and return to the main loop

*
*
*
* Main Memory TO RAM
*
          NAM      10      MMTORAM give an alpha prompt
*
          FCB      $000      indicate nonprogrammable
*
* initialize for the main loop
*

```

| | | | |
|--|---------|--------|--|
| | GOSUB | INMMTR | go do compile |
| | GOSUB | PATCH | get start address and say "LOADING PGM" |
| | CLRF 10 | | indicate we are in MM |
| | RCR 8 | | |
| | A=C | R< | set up end for CPGMHD |
| | ?NC XQ | CPGMHD | |
| | ?NC XQ | PUTPCF | set PC to begin of MM program |
| | C=0 | X | |
| | RAMSLCT | | select the status registers |
| | READ A | | get begin address of storage |
| | RCR 11 | | rotate to "C[6-3]" |
| | WRIT A | | begin address to reg A |
| | C=C+1 | M | skip the two header words |
| | M=C | | save code address-1 in "M" |
| * | | | |
| * start of the main loop of MMTORAM | | | |
| * | | | |
| DINSTR | A<>C | R< | |
| | N=C | | |
| | A<>C | R< | save begin of program in "N" |
| | ?NC XQ | NXLSST | end of this instruction to "A[3-0]" |
| | CLRF 9 | | clear the copy of the end flag |
| | ?FSET 6 | | have we found an END |
| | JNC+ | NOEND | |
| | SETF 9 | | copy status of flag 6 into flag 9 |
| NOEND | A<>C | R< | |
| | C<>N | | |
| | A<>C | R< | begin of instr. to "A[3-0]", end to "N[3-0]" |
| | SETF 8 | | indicate we do first byte of instr. |
| * | | | |
| * start of the 'instruction loop'. in this loop one instruction is saved | | | |
| * in ram memory | | | |
| * | | | |
| DOBYTE | ?NC XQ | NXBYTA | get a byte of this instr. |
| | A<>B | R< | current PC to "B[3-0]" |
| | R= 0 | | |
| | G=C | | fetch byte to "G" |
| | ?FSET 8 | | is this the first byte |
| | JNC+ | NOTB-E | no, skip byte test |
| | ?FSET 9 | | have we hit an end |
| | JC+ | NOTB-E | yes, skip byte test |
| | R= 1 | | |
| | A=C | X | save byte in "A[1-0]" |
| | C=C+1 | R | test for F byte |
| | JC+ | NOTB-E | skip B-E handling |
| | LDI | \$0BE | set up comparing constant |
| | ?A<C | R | if smaller than B, we can skip B-E handling |
| | JC+ | NOTB-E | |
| | GOSUB | DO B-E | take care of B-E bytes |
| NOTB-E | C=M | | get storage address-1 |
| | C=C+1 | M | |
| | C=0 | X | clear data space |
| | M=C | | save updated storage address |
| | R= 2 | | |
| | ?FSET 8 | | are we handling the first byte |
| | JNC+ | NOTONE | if not, skip the flagging of byte one |
| | LD@R 1 | | flag first byte |


```

NOTONE R= 0
      C=G          get byte
      CLRf 8       say first byte is done
      WRIT        put byte in ram
      C=N         get end of instr.
      R= 3
      A<>B      R<    address of byte to "C[3-0]" end to "A[3-0]"
      ?FSET 9     are we handling an END
      JC+        DOEND skip PC compare
      ?A#C      R<    have we handled the whole instr.
      JC-        DOBYTE no, do the next byte of this instr.
      JNC-       DINSTR yes, do next instr.

*
* make the END instr. in ram and test if it has to be a private END or not
* and if so make it private
*
DOEND  ?NC XQ      NXBYTA get next byte
      ST=C        save byte in "ST"
      C=M
      C=C+1      M    make new storage address
      C=ST       get byte back
      C=0        XS   clear first byte indicator
      WRIT       put byte in ram
      C=C+1      M    point to last address of program in ram
      M=C        save last used address for getting out first
*          free byte after program for the user
      LDI        $22F set up last byte of END
      C<>B      A    address and byte to "B"
      C=0        X
      RAMSLCT    select the status registers
      READ D
      RCR 13
      ST=C       user flags 0-3 into CPU 3-0
      C=B        A    get back address and byte
      R= 1
      ?FSET 2    must it be a private END
      JNC+       NOPRIV no
      LD@R 6     yes, indicate write protected file with $26F
NOPRIV WRIT      put byte in ram
*
* construct the header words for a ROM program file the first header word
* contains the length of the program in registers and the second word
* holds the status of the file ( private or not ) in the bit 8 and the
* number of bytes that are not hold by a completely filled up register in
* digit 1
*
      C=C-1      M    set up "C" for byte count
      A<>C      M    save end of program in "A"
      READ A     get start address
      A=A-C      M    number of bytes to "A(M)"
      C=0        A    set up for divide by 7 for header words
      R= 3
      LD@R 7
DIVIDE C=C+1      X    open at least 1 register
      ?A<C      M    end of division
      JC+        EOFDIV yes
      ?A#C      M    integer division
      JNC+       EOFDIV yes, exit
      A=A-C      M    subtract ( this is a dumb form of division )
      JNC-       DIVIDE

```

| | | | |
|--------|---------|--------|--|
| EOFDIV | A<>C | X | save number of regs in "A[2-0]" |
| | READ A | | get address of first header word |
| | A<>C | X | |
| | WRIT | | write first header word in ram |
| | C=C+1 | M | point to second header word |
| | A<>C | M | number of bytes to "C(M)" address to "A(M)" |
| | RCR 2 | | |
| | A<>C | X | bytecount to "A[1]" |
| | LDI | \$300 | say status is private |
| | ?FSET 3 | | does the user want it to be private |
| | JC+ | PRIVAT | yes |
| | LD@R 2 | | no, make it an open file |
| PRIVAT | A<>C | XS | combine status and number of bytes in "A(X)" |
| | A<>C | A | |
| | WRIT | | put second header word in ram |

*
* last thing to do is to compute the link in the first alpha label to the
* label that will be in front of the program, when it will be copied to MM
* for this we need to know how many empty bytes are left in the first reg
* and the distance from the program start to the first alpha label
*

| | | | |
|--|-------|-------|---|
| | A=C | M | start address-1 to "A(M)" |
| | C=0 | M | |
| | C=0 | XS | clear file status, preserve bytes |
| | RCR 1 | | |
| | A<>C | X | number of bytes to "A[0]" |
| | LDI | \$007 | |
| | A<>C | X | |
| | A=A-C | X | calculate 7-bytes (zero bytes in first reg) |
| | A=A+1 | X | |
| | A=A+1 | X | add displacement to previous link-1 |
| | B=A | A | save start-1 and offset |

*
* compute the distance to the first label starting from the first byte
*

| | | | |
|--------|-------|--------|---|
| | LDI | \$1CD | |
| | A<>C | X | comparing constant for G labels to "A(X)" |
| | C=B | M | get start-1 |
| | R= 1 | | |
| NXTBYT | C=C+1 | M | |
| | FETCH | | get next byte of the new file |
| | ?C#0 | XS | is this a start of an instr. |
| | JNC- | NXTBYT | no, continue with next byte |
| | ?A#C | R | is it a C byte |
| | JC- | NXTBYT | no, continue with the next byte |
| | ?A<C | X | one of the exceptions in C row |
| | JC- | NXTBYT | yes, look up another byte |
| | A<>C | M | start-1 to "C(M)", end to "C(M)" |
| | B=A | M | save end (bytes to change) |
| | A=A-C | M | number of bytes to first label in "C(M)" |

*
* we have the number of bytes to the previous link
*

| | | | |
|--|--------|---|--------------------------------------|
| | C=0 | A | set up for link distance computation |
| | C=B | X | |
| | RCR 11 | | offset value to "C(M)" |
| | A=A+C | M | total link distance in "A(M)" |
| | C=0 | A | set up for division |
| | R= 3 | | |
| | LD@R 7 | | |

```

DIVLNK  ?A<C      M      division finished
        JC+        LFOUND  yes, exit
        A=A-C      M
        C=C+1      X      increment register count
        JNC-        DIVLNK
LFOUND  A<>C      A      bytes in "C(M)", regs in "A(X)"
        C=C+C      M      double byte count
        C=0        X      clear dat area
        RCR 1      get bytes in place
        C=C+A      X      add register count
        ST=C        save second byte
        RCR 2      create first byte
        LD@R 1
        LD@R C
        C=B        M      get address for first byte
        WRIT        put it to ram
        C=C+1      M      point to second byte
        C=0        X
        C=ST        get second byte
        WRIT        put it in ram
EXIT    C=M        get back last word of program
        C=C+1      M      make first free word after program
        RCR 4
        C=0        M      leave only the address
        RCR 13      make the output the same as the input format
*
* now we initialize for DECOD and put the next free byte address back to
* the alpha register. then we exit through READY of compile, saying "READY"
* and giving an audio warning
*
        C<>B      A
        A=0        A
        C=0        A
        SETF 8
        CLRF 9
        WRIT 0
        WRIT P      initialization done
        GOSUB      DECOD  put address in alpha
*
* we can exit through READY when we make sure that flag 0 of reg E is
* cleared.
*
EXIT2   READ E
        RCR 3
        ST=0
        C=ST        indicate we have finished compile
        RCR 11
        WRIT E
        GOTO        READY  and do exit and say "READY"
*
*
*
* ROMSUM
*
        NAM        ROMSUM
*

```

| | | |
|----------|-------|---|
| GOSUB | CODE | get the page to romsum |
| RCR 11 | | P to "C[3]" |
| C=0 | X | |
| C=-C-1 | X | |
| RCR 11 | | PFFF to "C(M)" |
| C=0 | X | |
| WRIT | | clear the romsum word |
| R= 5 | | |
| C=0 | R< | |
| C<>B | A | start address to "B(M)" |
| C=0 | A | |
| RAMSLCT | | select the status registers |
| PRPHSLCT | | and deselect all peripherals |
| LDI | \$300 | |
| C=C+C | X | |
| C=C+C | X | |
| A=C | A | "A(X)" has %1100 0000 0000 |
| C=C+1 | X | |
| C<>B | A | "B(X)" has %1100 0000 0001 "C(M)" has start |

*
* the checksum is computed by adding all the words to each other and when a
* carry occurs to add this carry to the total checksum. the final checksum
* is made by taking the 2's complement of the added words.
*
* detection of the carry is done by making the two left bits high. when the
* carry in the right ten bits occur, it will be transferred to bit 12 and
* this is the carry bit. we add the carry and restore the status of the two
* left bits in the S&X field of "A" by adding %1100 0000 0001 to "A(X)"
*

| | | | |
|------|-------|-------|---|
| ROMS | FETCH | | get a rom word |
| | A=A+C | X | and add it to the checksum |
| | JNC+ | 2 | we didn't had a carry |
| | A=A+B | X | add one to the checksum and restore bit 10 + 11 |
| | C=0 | X | |
| | C=C+1 | M | point to next word to add |
| | ?C#0 | R< | have we had the whole rom yet |
| | JC- | ROMS | no, do next word |
| | A=A-B | X | decrement "A(X)" and get rid of bit 10 and 11 |
| | LDI | \$3FF | |
| | A<>C | A | |
| | C=A-C | A | get the 2's complement of bit 0-9 to "C(X)" |
| | C=C-1 | M | point to checksum word |
| | WRIT | | put checksum into the ram |
| | RTN | | |

*
*
*
* DISASSEMBle
*

| | | |
|--|--------|--------|
| | NAM | DISASM |
| | SETHEX | |
| | ?NC XQ | CLA |
| | LDI | \$006 |
| | RCR 11 | |
| | LDI | \$009 |
| | A=C | A |
| | READ X | |
| | C=C+1 | A |
| | WRIT X | |

\$006009 to "A[5-0]"
get word to disassemble
save address of next word for the user

```

C=C-1      A
RCR 11
FETCH      get the word to disassemble
WRIT Y     word and address to reg Y for MNEM
RCR 7      left justify string for conversion to ascii
DISAS1 LDI      $003    load first ascii digit of a character
RCR 13
R= 0
?A<C      R      is it a number 9-0
JNC+      5      yes, then we are ready with converting
*
* when the hex digit is A-F we have to make it 1-6 and add 1 to the first
* digit of the ascii character
*
C=A-C      R
C=-C      R      convert digit to 1-6
R= 1
C=C+1      R      make first digit 4
C<>B      A      save string with converted character in "B"
READ M
RCR 12
C<>B      X
C<>B      XS      put new character in place
WRIT M     save converted digits
C<>B      A      get string back for next character
A=A-1      M      are we finished with characters
JNC-      DISAS1    no, do one more
*
* the disassembled word and address are written in alpha as follows
* AAAA_DDD_C_ where '_' stands for a space the C stands for the ascii
* representation of the word and is added in the conversion part
*
READ M     get string back
A=C      A      and save it in "A"
R= 5
C=0      R<      leave only the address part
RCR 6      right justify it
WRIT N     and write it to reg N of the alpha register
A<>C      A      get string back
RCR 8
LD@R 2     load three spaces after the data word
LD@R 0
LD@R 2
LD@R 0
LD@R 2
LD@R 0
R= 13
LD@R 2     load one space before data word
LD@R 0
WRIT M     write string to reg M of the aloha register
READ Y
*
* get the data word back and decide by which character it has to be
* represented. the upper bits of the word are ignored for the ascii
* character as is bit three of the second digit ( bit 7 )
*
* words where the second digit is 4,5,6 or 7 are represented by the
* characters which are picked from the ASCTBL at $2CF0. which character
* is used is determined by the first digit of the word
*

```

* words where the second digit is 2 or 3 are represented by their ascii
 * value
 *
 * words where the second digit is 0 or 1 are represented by the characters
 * that are represented by the value of the word added with \$40
 *

| | | |
|----------|-------|--|
| ST=C | | |
| CLRF 7 | | |
| C=ST | | bit 7 of the word is ignored |
| ?FSET 6 | | is first digit 4,5,6 or 7 |
| JNC+ | 8 | no, test for digit 0,1 or 2,3 |
| R= 3 | | load address of ASCTBL the character to |
| LD@R 2 | | be picked is determined by the first digit |
| LD@R C | | of the word we use |
| LD@R 0 | | |
| RCR 11 | | address to "C(M)" |
| FETCH | | get ascii value |
| JNC+ | 5 | end of conversion |
| ?FSET 5 | | is second digit a 2 or 3 |
| JC+ | 3 | yes, represent by ascii #20 upto \$3F |
| SETF 6 | | no, represent by ascii \$40 upto \$5F |
| C=ST | | get converted word \$00 upto #1F |
| A=C | A | save character in "A[1-0]" |
| READ M | | |
| RCR 2 | | |
| A<>C | X | |
| A<>C | XS | |
| RCR 12 | | append character to alpha string |
| WRIT M | | |
| ?FSET 13 | | is there a program running |
| ?C RTN | | yes, return to mainframe |
| ?NC GO | AVIEW | no, put string in display |

*
 *
 *

* COPY Rom

*

| | | |
|--------|-------|--|
| NAM | COPYR | |
| GOSUB | COD | get SD to "C[1-0]" |
| R= 6 | | |
| RCR 9 | | S to "C[6]" |
| A=0 | A | |
| A<>C | R | leave S000 in "A(M)" |
| RCR 13 | | |
| B=0 | A | |
| C<>B | R | D to "B[6]" |
| A<>C | R | |
| A=C | R | S to "C[6]" |
| A=A+1 | R | S+1 to "A[6]" for stop criterium |
| COPY | FETCH | get a source word |
| C<>B | R | get destination address |
| WRIT | | write source word to destination |
| C<>B | R | get source address back |
| C=C+1 | M | increment source and destination address |
| ?A#C | R | have we had the whole page yet |
| JC- | COPY | no, do the next word |
| LDI | \$027 | |
| ?NC GO | TONEB | give a beep and return to mainframe |

*
 *

```

*
* ROMCHecKX
*
*      NAM      ROMCHKX
*
      C=0      A
      RAMSLCT
      READ X
      ?NC XQ      BCDBIN
      C<>B      X      xrom number to check to "B(X)"
      C=0      A
      R= 6
      LD@R 4
      C<>B      M
      CLRf 9      reset error flag
TRYNRM C=B      M
      R= 6
      C=C+1      R      point to next rom
      JNC+      4      rom not found yet
      GOTO      NOROM  say specified rom not found
      C<>B      M
      C=B      M
      FETCH      get rom ID
      ?C#0      X      is this an existent rom
      JNC-      TRYNRM no, try next rom
      A<>B      X
      B=A      X
      ?A#C      X      is it the one we want to check
      JC-      TRYNRM no, do next rom
      ?NC XQ      CLLCDE
      A=0      MS
      ?NC XQ      GENNUM  get xrom number in 2 characters to display
      LDI      $020
      WRIT E      append a blank
      C=0      A
      R= 5
      LD@R F
      LD@R F
      LD@R E
      A=C      M
      C=B      M
      C=C+A      M      construct address of revision number
      FETCH
      WRIT E      append one rev. character to display
      C=C-1      M
      FETCH
      WRIT E      append second rev. character to display
      LDI      $02D
      WRIT E      seperate rev. level and characters with a "-"
      C=C-1      M
      FETCH
      WRIT E      append one revision level to display
      C=C-1      M
      FETCH
      WRIT E      append second revision level to display
      ?NC XQ      MESSL
      RMB      $04      " TST"
*
* display has now message NN RR-LL TST where NN is the desired xrom number
* RR is the revision code and LL is the revision level.
*

```

```

C=B      A      get start address of rom to "C(M)"
LDI      $300
C=C+C    X
C=C+C    X
A=C      A      "A(X)" = 1100 0000 0000 start address to "A(M)"
C=C+1    X
C<>B     X      "B(X)" = 1100 0000 0001
C=0      A
C=C+1    A
RCR 8
A<>C     M      "C(M)" has start address
A=A+C    M      "A(M)" has end address
*
* in this loop all the rom words are added to create the chcksum. when the
* rom is alright the checksum should end up as zero.
*
WORD      FETCH      get word
C=C+1     M          increment address
A=A+C     X          add to checksum
JNC+      2
A=A+B     X          add cary to checksum and fix high order bits
*
* we need to fix the high order bits to detect the next carry.
*
?A#C      M          are we finished yet
JC-        WORD      no, do next word
A=A-B     X
A=A-1     X          is checksum zero
JC+        2
SETF 9          no, set error flag
READ D
READ D
READ D          get the TST characters of the display
?FSET 9        is rom ok
JC+           BAD     no, append message BAD
?NC XQ        MESSL    yes, append message OK
RMB          $03       "OK "
JNC+          EXRCHK
?NC XQ        MESSL
RMB          $03       "BAD"
JNC+          EXRCHK
NOROM SETF 9        set error flag
?NC XQ        CLLCDE
?NC XQ        MESSL
RMB          $07       "NO ROM "
A<>B     X
A=0       MS
?NC XQ    GENNUM
?NC XQ    LEFTJ
EXRCHK ?NC XQ    ENCP00
?NC XQ    STMSGF
?FSET 9
?NC RTN
SETF 7
?NC GO    SKP
*
*
*
* SYNTesize
*

```



```

      NAM          SYNT
*
      FCB          $000    say that this function is non programmable
      READ Y       get the prefix byte
      ?NC XQ       CHK#S    check for alpha data
      SETHEX
      ?NC XQ       BCDBIN
      C<>B         X       save prefix in "B[1-0]"
      READ X       get the postfix byte
      ?NC XQ       CHK#S    check for alpha data
      SETHEX
      ?NC XQ       BCDBIN
      RCR 2        postfix to "C[13-12]"
      C<>B         X
      RCR 2        postfix to "C[11-10]" prefix to "C[13-12]"
      ?NC GO       INSLIN   insert the byte in main memory
*
*
*
* Go to End
*
      NAM          GE
*
      C=0          A
      WRIT A
      WRIT B       delete all pending return addresses
      CLR 10       say PC is in RAM
      READ C       get address of register of .END
      R= 3
      LD@R 6       make address of .END
      R=3
      A<>C         R<      initialize for PUTPC
      ?NC GO       PUTPCF   place PC at .END and set line number to $FFF
*
*
*
* Initialize Main Memory To Ram
*
* here we save the return address for MMTORAM in the first 4 digits of
* scratch register P. these digits are used by COMPILE to get it's return
* address from, in case we use it as a subroutine.
*
      INMMTR POP    get return address
      RCR 4
      C=0          M
      RCR 3        leave return address only
      WRIT P
*
* second thing we have to do here is to find out whether the user wants to
* delete the labels from the program, or just wants to do a normal compile
* and load program. this is indicated by flag 3 of the user flags.
*

```

```

READ D
RCR 13
ST=C                                type of compile to flag 0
CLRf 1                             delete status of user flag 2
SETF 2                             say we use compile as a subroutine
READ E
RCR 3
C=ST                                put flags in place
RCR 11
WRIT E                             store flags in reg E
GOTO                                COMPIL
*
*
*
* PATCH MMTORAM
*
PATCH  GOSUB          COD      get the start address
        WRIT A
        ?NC XQ        CLLCDE
        ?NC XQ        MESSL
        RMB           $OC      "LOADING PGM "
        ?NC XQ        LEFTJ    put message in display
        ?NC XQ        ENCP00
        GOTO          FPAL     find links and return to MMTORAM
*
*
*
* ROM NAME
*
        NAM           ESMLDL-OS
*
        RTN
*
*
*
* Make Private
*
        NAM           10 MKPR
*
        FCB           $000
*
        GOSUB          FPAL
        RCR 8
        A<>C          R<
        ?NC XQ        OFEE     get third byte to "ST" from address in "A[3-0]"
        CLRf 5
        SETF 6         set private bit
        C<>ST
        SETF 12        set indication for cpu of private
        ?NC XQ        PTBYTA
        ?NC GO        NFRPU    return and update private status
*
*

```

```

*
* no entry points used
*
*
*      ORG      $FFB
*
*      FCB      $001      revision level is  A
*      FCB      $020
*
*      FCB      $013      revision code is AS
*      FCB      $001
*
*      FCB      $168      checksum
*
*
*      END
*
*****
*
*                               END OF ASSEMBLY LISTING
*
*****
*
*
*
*                               ERAMCO SYSTEMS
*
*
*

```